

# Recursive functions

recursive := a function which calls itself.

break a problem up into smaller/simpler versions of the problem, solve those problems recursively, combine results to solve original problem.

## Example

$$n! = n(n-1)(n-2) \dots 1$$

$$0! = 1$$

function factorial(n)

...

if  $n = 0$ , return 1 // base case  
else, return  $n * \text{factorial}(n-1)$  // recursive step

$$n! = n[(n-1)!]$$

↑ smaller version of the problem

iterative version

function factorial2(n)

$$f = 1$$

for  $i = 1$  to  $n$

$$f = f * i$$

Return  $f$  // ~~1 \* 2 \* 3 \* ... \* n = n!~~

Example  $a^n = a \cdot a^{n-1}$

function  $\text{expon}(a, n)$

if  $n=0$ , return 1 // base case

else, return  $a \cdot \underbrace{\text{expon}(a, n-1)}_{a^{n-1}}$

fast exponentiation

$$a^n = \begin{cases} (a^{\frac{n}{2}})^2 & \text{when } n \text{ is even} \\ a(a^{\frac{n-1}{2}})^2 & \text{when } n \text{ is odd} \end{cases}$$

function  $\text{fast\_expon}(a, n)$

if  $n=0$ , return 1

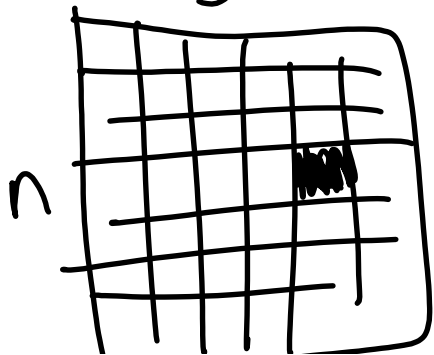
if  $n=1$ , return  $a$

if  $n$  is even,

return  $(\text{fast\_expon}(a, \frac{n}{2}))^2$   
 else  
 return  $a \cdot (\text{fast\_expon}(a, \frac{n}{2}))^2$

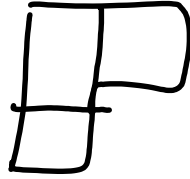
$$\begin{aligned}
 T(n) &= c + T\left(\frac{n}{2}\right) \\
 &= c + (c + T\left(\frac{n}{4}\right)) \\
 &\vdots \\
 &= O(\log n)
 \end{aligned}$$

Tiling with triminoes

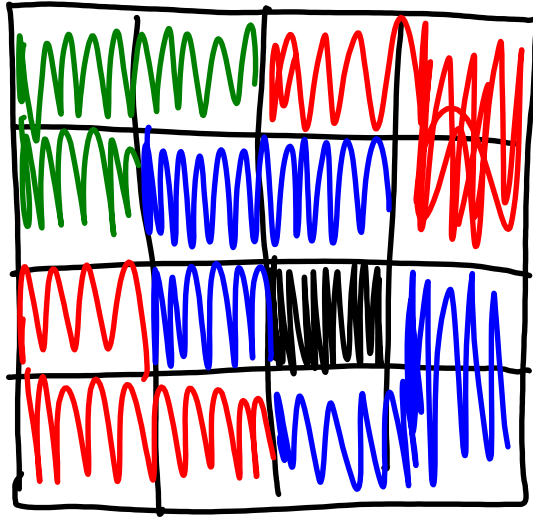


$n \times n$  checkerboard w/  
 $n = 2^k$ ,  $k \in \mathbb{Z}^+$  1 cell  
 missing

$n$



triminoes



proc