

# CSCE 222-200 Discrete Structures for Computing

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**Instructor:** Dr. Jianer Chen

**Office:** PETR 428

**Phone:** (979) 845-4259

**Email:** chen@cse.tamu.edu

**Office Hours:** T+R 2:00pm–3:30pm

**Teaching Assistant:** Evan Kostov

**Office:** EABC Cubicle 6

**Phone:** (469) 996-5494

**Email:** evankostov@tamu.edu

**Office Hours:** MW 4:00pm–5:00pm

## Assignment #5 Solutions

1. Find a big- $O$  estimate for the function  $f(n)$  that satisfies the recurrence relation

$$f(2) = 1; \quad \text{and} \quad \text{for } n > 2, f(n) = 2f(\sqrt{n}) + \log_2 n.$$

(*Hint:* Make the substitution  $m = \log_2 n$ . You can assume that  $m$  is a power of 2.)

**Solution.** Let  $m = \log_2 n$ , i.e.,  $n = 2^m$ . Define  $F(m) = f(n) = f(2^m)$ . We have

$$F(1) = f(2^1) = f(2) = 1, \text{ and}$$

for  $m > 1$ ,

$$\begin{aligned} F(m) = f(2^m) = f(n) &= 2f(\sqrt{n}) + \log_2 n = 2f(\sqrt{2^m}) + m = 2f(2^{m/2}) + m \\ &= 2F(m/2) + m. \end{aligned}$$

Solving  $F(1) = 1$  and for  $m > 1$ ,  $F(m) = 2F(m/2) + m$  is not difficult, as we have done similar recurrence relations. For a general  $k \leq m$ , we can verify

$$F(m) = 2^k F(m/2^k) + k \cdot m.$$

Letting  $k = \log_2 m$  gives

$$F(m) = m \cdot F(1) + \log_2 m \cdot m = m + m \log_2 m.$$

Recall that  $m = \log_2 n$  and  $n = 2^m$ . We have, for a general  $n \geq 2$ ,

$$\begin{aligned} f(n) = f(2^m) = F(m) &= m + m \log_2 m = \log_2 n + \log_2 n \cdot \log_2 \log_2 n \\ &= O(\log_2 n \cdot \log_2 \log_2 n). \end{aligned}$$

2. A coin is flipped  $n$  times where each flip comes up either heads or tails. How many possible outcomes (assuming that  $n$  is even and that  $k \leq n$ )

- are there in total?
- contain exactly  $k$  heads?
- contain at least  $k$  heads?
- contain the same number of heads and tails?

Give an explanation to your answer to each of the questions.

**Solution.** If we treat head  $H$  as the bit 0 and tail  $T$  as the bit 1, then each outcome of  $n$  coin-flippings uniquely corresponds to a binary string of length  $n$ . This gives a one-to-one correspondence between the set of outcomes of  $n$  coin-flippings and the set of binary strings of length  $n$ .

(a) As we know, the total number of binary strings of length  $n$  is  $2^n$ . Thus, there are totally  $2^n$  possible outcomes for  $n$  coin-flippings.

(b) Each outcome with exactly  $k$  heads corresponds uniquely to  $k$  positions in the binary string of length  $n$  at which the bit is 0 (and all other bits are 1). Since there are  $\binom{n}{k}$  different ways to pick  $k$  positions in a binary string of length  $n$ , the total number of outcomes that contain exactly  $k$  heads is equal to  $\binom{n}{k}$ .

(c) As given in (b), for each  $h$ ,  $k \leq h \leq n$ , the number of outcomes that contain exactly  $h$  heads is equal to  $\binom{n}{h}$ . The total number of outcomes that contain at least  $k$  heads is equal to the sum of the numbers of outcomes that contain exactly  $h$  heads for all  $k \leq h \leq n$ . That is, the total number of outcomes that contain at least  $k$  heads is equal to  $\sum_{h=k}^n \binom{n}{h}$ .

(d) Since  $n$  is even,  $n/2$  is an integer. An outcome containing the same number of heads and tails contains exactly  $n/2$  heads. As a consequence of (b), the total number of outcomes that contain the same number of heads and tails is equal  $\binom{n}{n/2}$ .

**3.** Suppose that  $k$  and  $n$  are integers with  $1 \leq k \leq n$ . Prove the *hexagon identity*:

$$\binom{n-1}{k-1} \binom{n}{k+1} \binom{n+1}{k} = \binom{n-1}{k} \binom{n}{k-1} \binom{n+1}{k+1}.$$

**Solution.** We prove the equality using the formula  $\binom{n}{k} = n!/(k! \cdot (n-k)!)$ .

$$\begin{aligned} \text{LHS} &= \binom{n-1}{k-1} \binom{n}{k+1} \binom{n+1}{k} \\ &= \frac{(n-1)!}{(n-k)!(k-1)!} \cdot \frac{n!}{(n-k-1)!(k+1)!} \cdot \frac{(n+1)!}{(n-k+1)!k!} \\ &= \frac{(n-1)! \cdot n! \cdot (n+1)!}{(k-1)! \cdot k! \cdot (k+1)! \cdot (n-k-1)! \cdot (n-k)! \cdot (n-k+1)!} \end{aligned}$$

and

$$\begin{aligned} \text{RHS} &= \binom{n-1}{k} \binom{n}{k-1} \binom{n+1}{k+1} \\ &= \frac{(n-1)!}{(n-k-1)!k!} \cdot \frac{n!}{(n-k+1)!(k-1)!} \cdot \frac{(n+1)!}{(n-k)!(k+1)!} \\ &= \frac{(n-1)! \cdot n! \cdot (n+1)!}{(k-1)! \cdot k! \cdot (k+1)! \cdot (n-k-1)! \cdot (n-k)! \cdot (n-k+1)!} \end{aligned}$$

Thus, we have LHS = RHS. The equality is proved.

4. Prove that if  $E$  and  $F$  are independent events, then  $\overline{E}$  and  $\overline{F}$  are also independent events.

**Proof.** Recall that  $E$  and  $F$  are independent means that  $\Pr[E \cap F] = \Pr[E] \cdot \Pr[F]$ . Thus, the question asks to prove  $\Pr[\overline{E} \cap \overline{F}] = \Pr[\overline{E}] \cdot \Pr[\overline{F}]$ . We have

$$\begin{aligned} \Pr[\overline{E} \cap \overline{F}] &= \Pr[\overline{E \cup F}] \\ &= 1 - \Pr[E \cup F] \\ &= 1 - (\Pr[E] + \Pr[F] - \Pr[E \cap F]) \\ &= 1 - (\Pr[E] + \Pr[F] - \Pr[E] \cdot \Pr[F]) \\ &= 1 - \Pr[E] - \Pr[F] + \Pr[E] \cdot \Pr[F] \\ &= (1 - \Pr[E]) \cdot (1 - \Pr[F]) \\ &= \Pr[\overline{E}] \cdot \Pr[\overline{F}]. \end{aligned}$$

The first equality used De Morgan's law (note that  $E$  and  $F$  are sets), the second and seventh equalities used the formula for the probability of event complements, the third equality used the formula for the probability of union of non-disjoint events, and the fourth equality used the given condition that  $E$  and  $F$  are independent events.

This proves that the events  $\overline{E}$  and  $\overline{F}$  are also independent.

5. Suppose that we roll a fair die until a 6 comes up.

- What is the probability that a 6 comes up in our  $n$ -th rolling?
- What is the expected number of times we roll the die? (*Hint:* You need to find the value for the sum  $1 + 2(5/6) + 3(5/6)^2 + \dots + k(5/6)^{k-1} + \dots$ .)

**Solution.** Note that when rolling a fair die, the probability that 6 comes up is  $1/6$ , and the probability that 6 does not come up is  $5/6$ .

(a) According to the question statement, the game will stop when 6 comes up. Thus, that 6 comes up in the  $n$ -th rolling implies that 6 did not come up in the first  $(n - 1)$ st rollings. Thus, the probability that a 6 comes up in the  $n$ -th rolling is  $(5/6)^{n-1}(1/6)$ .

(b) For each integer  $n \geq 1$ , define a random variable  $X_n$  such that  $X_n = n$  if a 6 comes up in the  $n$ -th rolling, and  $X_n = 0$  otherwise. By (a), the probability that  $X_n = n$  is equal to  $(5/6)^{n-1}(1/6)$ , so the probability that  $X_n = 0$  is  $1 - (5/6)^{n-1}(1/6)$ . Therefore, the expected value of the random variable  $X_n$  is

$$\begin{aligned} \mathbf{Ex}[X_n] &= n \cdot \Pr[X_n = n] + 0 \cdot \Pr[X_n = 0] \\ &= n \cdot (5/6)^{n-1}(1/6) + 0 \cdot (1 - (5/6)^{n-1}(1/6)) \\ &= n \cdot (5/6)^{n-1}(1/6). \end{aligned}$$

Define a new random variable  $X = \sum_{n=1}^{\infty} X_n$ . Since in every case, there is at most one  $X_n$  that is not equal 0 (i.e., when a 6 comes up in the  $n$ -th roll, we have  $X_n = n$ , and for all  $i \neq n$ ,  $X_i = 0$ ),  $X$  is the number of times we roll the die to have a 6 comes up. Therefore, question (b) is asking the expected value of  $X$ . We have

$$\begin{aligned} \mathbf{E}\mathbf{x}[X] &= \mathbf{E}\mathbf{x}\left[X = \sum_{n=1}^{\infty} X_n\right] = \sum_{n=1}^{\infty} \mathbf{E}\mathbf{x}[X_n] \\ &= \sum_{n=1}^{\infty} n \cdot (5/6)^{n-1} (1/6) = (1/6) \sum_{n=1}^{\infty} n \cdot (5/6)^{n-1}, \end{aligned} \quad (1)$$

where the second equality has used the linearity of expectations.

To get the final value of  $\mathbf{E}\mathbf{x}[X]$ , we need to find the value of

$$\sum_{n=1}^{\infty} n \cdot (5/6)^{n-1} = 1 + 2(5/6) + 3(5/6)^2 + \cdots + n(5/6)^{n-1} + \cdots.$$

For this, we have

$$\begin{aligned} \sum_{n=1}^{\infty} n \cdot (5/6)^{n-1} &= 1 + 2(5/6) + 3(5/6)^2 + 4(5/6)^3 + \cdots + n(5/6)^{n-1} + \cdots \\ &= 1 + (5/6) + (5/6)^2 + (5/6)^3 + \cdots + (5/6)^{n-1} + \cdots \\ &\quad + (5/6) + 2(5/6)^2 + 3(5/6)^3 + \cdots + (n-1)(5/6)^{n-1} + \cdots \\ &= 1 + (5/6) + (5/6)^2 + (5/6)^3 + \cdots + (5/6)^{n-1} + \cdots \\ &\quad + (5/6) \left( \sum_{n=1}^{\infty} n \cdot (5/6)^{n-1} \right) \end{aligned} \quad (2)$$

By the formula for the summation of geometric sequences, we have

$$1 + (5/6) + (5/6)^2 + (5/6)^3 + \cdots + (5/6)^{n-1} + \cdots = 1/(1 - (5/6)) = 6.$$

Therefore, from the above equations in (2), we have

$$\sum_{n=1}^{\infty} n \cdot (5/6)^{n-1} = 6 + (5/6) \sum_{n=1}^{\infty} n \cdot (5/6)^{n-1}.$$

Solving this for  $\sum_{n=1}^{\infty} n \cdot (5/6)^{n-1}$ , we get

$$\sum_{n=1}^{\infty} n \cdot (5/6)^{n-1} = 36.$$

Bringing this into (1) gives

$$\mathbf{E}\mathbf{x}[X] = (1/6) \sum_{n=1}^{\infty} n \cdot (5/6)^{n-1} = (1/6) \cdot 36 = 6.$$