

CSCE 222-200 Discrete Structures for Computing

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Assignment #1 Solutions

1. Prove by induction that $3^n < n!$ if n is an integer greater than 6.

Proof. We prove the inequality by induction on the integer $n > 6$.

Basis Step. Since $n > 6$, the basis case is $n = 7$. In this case, we have LHS = $3^7 = 2187$, and RHS = $7! = 5040$. Thus, LHS < RHS. The inequality holds true.

Inductive Step. Assume inductively that the inequality holds true for all n , when $7 \leq n \leq k$.

Now consider the inequality for $n = k + 1$. We have

$$\text{LHS} = 3^n = 3^{k+1} = 3^k \cdot 3 < k! \cdot 3 < k! \cdot (k + 1) = (k + 1)! = n! = \text{RHS},$$

where in the first inequality, we have used the inductive hypothesis $3^k < k!$, while in the second inequality, we have used the fact that $3 < k + 1$ (because $7 \leq k$). As a result, for $n = k + 1$, we still have the inequality $3^n < n!$ hold true.

By mathematical induction, this proves that $3^n < n!$ holds true for all $n > 6$.

2. (a) Find a formula for $\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \cdots + \frac{1}{n(n+1)}$ by examining the values of this expression for small values of n .

(b) Prove by induction the formula you conjectured in part (a).

Proof. (a) You may try the formula in any way you like. On the other hand, if you observe that $\frac{1}{n(n+1)} = \frac{1}{n} - \frac{1}{n+1}$ (for all $n \geq 1$), then the formula of the given summation becomes obvious:

$$\begin{aligned} & \frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \cdots + \frac{1}{n(n+1)} \\ &= \left(\frac{1}{1} - \frac{1}{2} \right) + \left(\frac{1}{2} - \frac{1}{3} \right) + \cdots + \left(\frac{1}{n} - \frac{1}{n+1} \right) \\ &= 1 - \frac{1}{n+1}. \end{aligned}$$

(b) Now we verify the above guessed equality

$$\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \cdots + \frac{1}{n(n+1)} = 1 - \frac{1}{n+1}$$

by induction on the integer $n \geq 1$.

Basis Step. For the basis case $n = 1$, we have LHS = $\frac{1}{1 \cdot 2} = \frac{1}{2}$, and RHS = $1 - \frac{1}{1+1} = \frac{1}{2}$. Thus, LHS = RHS. The equality holds true.

Inductive Step. Assume inductively that the equality holds true for all n , when $1 \leq n \leq k$.

Now consider the equality for $n = k + 1$. We have (note that $n \geq 2$ because $k \geq 1$)

$$\begin{aligned} \text{LHS} &= \frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \cdots + \frac{1}{(n-1)n} + \frac{1}{n(n+1)} \\ &= \frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \cdots + \frac{1}{k(k+1)} + \frac{1}{(k+1)(k+2)} \\ &= \left(1 - \frac{1}{k+1}\right) + \frac{1}{(k+1)(k+2)} \\ &= \left(1 - \frac{1}{k+1}\right) + \left(\frac{1}{(k+1)} - \frac{1}{(k+2)}\right) \\ &= 1 - \frac{1}{(k+2)} \\ &= 1 - \frac{1}{n+1} \\ &= \text{RHS}, \end{aligned}$$

where in the third equality, we have used the inductive hypothesis $\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \cdots + \frac{1}{n(n+1)} = 1 - \frac{1}{n+1}$ for $n = k$, and in the fourth equality, we have used the fact $\frac{1}{(k+1)(k+2)} = \frac{1}{k+1} - \frac{1}{k+2}$. As a result, for $n = k + 1$, we also have the equality $\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \cdots + \frac{1}{n(n+1)} = 1 - \frac{1}{n+1}$ hold true.

By mathematical induction, this proves that $\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \cdots + \frac{1}{n(n+1)} = 1 - \frac{1}{n+1}$ holds true for all $n \geq 1$.

3. Prove by contradiction that if x^3 is irrational, then x is irrational.

Proof. Assume the contrary that x is not irrational, i.e., x is a rational number. Then x can be written as the ratio of two integers m and n : $x = m/n$. Take the third power on both sides, we get $x^3 = (m/n)^3 = m^3/n^3$. Since m and n are integers, m^3 and n^3 are also integers. As a result, $x^3 = m^3/n^3$ would also be a rational number, contradicting the assumption that x^3 is irrational.

This contradiction proves that x must be irrational.

4. Prove by contradiction that $\sqrt{2} + \sqrt{3}$ is irrational.

Proof. Assume the contrary that $\sqrt{2} + \sqrt{3}$ is a rational number, i.e., $\sqrt{2} + \sqrt{3} = m/n$, where m and n are integers. Take square on both sides, we get

$$(\sqrt{2} + \sqrt{3})^2 = m^2/n^2.$$

Since $(\sqrt{2} + \sqrt{3})^2 = 2 + 2 \cdot \sqrt{2} \cdot \sqrt{3} + 3 = 5 + 2 \cdot \sqrt{6}$, we get

$$\sqrt{6} = (m^2/n^2 - 5)/2 = (m^2 - 5n^2)/(2n^2).$$

Since m and n are integers, $m^2 - 5n^2$ and $2n^2$ are also integers. As a result, $\sqrt{6} = (m^2 - 5n^2)/(2n^2)$ is a rational number, which can be written as

$$\sqrt{6} = r/q, \tag{1}$$

where r and q are integers with no common factor larger than 1. Square both sides of equality (1) then multiply both sides by q^2 , we get

$$6q^2 = r^2. \tag{2}$$

Since $6q^2$ is divisible by 2, r^2 is divisible by 2, which, as we discussed in class, leads to the conclusion that r is also divisible by 2. Write $r = 2k$, where k is an integer, then $r^2 = 4k^2$. Replacing r^2 in (2) by $4k^2$, then divide both sides by 2, we get

$$3q^2 = 2k^2. \tag{3}$$

However, this implies that q^2 is divisible by 2, which, again by our discussion in class, will further leads to the conclusion that q is divisible by 2. Thus, we would derive that both r and q are divisible by 2, contradicting our assumption that r and q have no common factor larger than 1.

This contradiction proves that $\sqrt{2} + \sqrt{3}$ is irrational.

5. Let p , q , and r be the propositions:

p : You get an A on the final exam;

q : You do every exercise in this book;

r : You get an A in this class.

Write each of the propositions below using p , q , r and logical operators:

- (a) You get an A in this class, but you do not do every exercise in this book.
- (b) You get an A on the final, you do every exercise in this book, and you get an A in this class.
- (c) To get an A in this class, it is necessary for you to get an A on the final.
- (d) You get an A on the final, but you do not do every exercise in this book; nevertheless, you get an A in this class.

- (e) Getting an A on the final and doing every exercise in this book is sufficient for getting an A in this class.
- (f) You will get an A in this class if and only if you either do every exercise in this book or you get an A on the final.

Solution.

- (a) $r \wedge \neg q$. (b) $p \wedge q \wedge r$. (c) $r \rightarrow p$.
- (d) $p \wedge \neg q \wedge r$. (e) $(p \wedge q) \rightarrow r$. (f) $r \leftrightarrow (p \vee q)$.

6. For each of the compound propositions below, use the conditional-disjunction equivalence to find an equivalent compound proposition that does not involve conditionals:

- (a) $\neg p \rightarrow \neg q$. (b) $(p \vee q) \rightarrow \neg p$. (c) $(p \rightarrow \neg q) \rightarrow (\neg p \rightarrow q)$.

Solution. The formula for the conditional-disjunction equivalence is $p \rightarrow q \equiv \neg p \vee q$.

- (a) $\neg p \rightarrow \neg q \equiv \neg(\neg p) \vee (\neg q) \equiv p \vee \neg q$.
- (b) $(p \vee q) \rightarrow \neg p \equiv \neg(p \vee q) \vee (\neg p) \equiv (\neg p \wedge \neg q) \vee (\neg p)$.
- (c) $(p \rightarrow \neg q) \rightarrow (\neg p \rightarrow q) \equiv (\neg p \vee \neg q) \rightarrow (p \vee q) \equiv \neg(\neg p \vee \neg q) \vee (p \vee q) \equiv$
 $\equiv (p \wedge q) \vee (p \vee q) \equiv (p \vee (p \vee q)) \wedge (q \vee (p \vee q)) \equiv (p \vee q) \wedge (p \vee q) \equiv p \vee q$.