

# Sparse Matrix Methods

## Chapter 5 lecture notes

### QR factorization

Tim Davis

2011

## Householder reflections

- $H = I - \beta vv^T$ ;  $\beta$  a scalar and  $v$  a vector
- can choose  $v$  and  $\beta$  based on  $x$  so that  $Hx$  is all zero except  $(Hx)_1 = \pm\|x\|_2$ .
- $Hx = x - v(\beta(v^T x))$
- $H$  is symmetric ( $H = H^T$ )
- $H$  is orthogonal  $HH^T = H^T H = I$
- if  $x_1$  is nonzero, then  $x$  and  $v$  have the *same* nonzero pattern ( $\mathcal{V} = \mathcal{X}$ ).
- if  $x_1$  is zero, then permute the rows of  $x$  to make it so

## QR factorization

- pick  $H_1$  based on  $A(:, 1)$ , to zero out all but  $A(1, 1)$
- then  $A = H_1 * A$
- repeat for 2nd column, zeroing out everything below the diagonal
- etc, until  $A$  becomes upper triangular

### Theorem (Golub)

The QR factorization  $QR = A$ , where  $A \in \mathbb{R}^{m \times n}$  and  $m \geq n$ , is  
 $Q = H_1 H_2 \cdots H_n = \prod_{k=1}^n H_k$  and  
 $R = Q^T A = H_n \cdots H_2 H_1 A = (\prod_{k=1}^n H_k) A = A^{[n]}.$

## MATLAB prototype: right-looking QR

```
function [V,Beta,R] = qr_right (A)
[m n] = size (A) ;
V = zeros (m,n) ;
Beta = zeros (1,n) ;
for k = 1:n
    [v,beta,s] = gallery ('house', A (k:m,k), 2) ;
    V (k:m,k) = v ;
    Beta (k) = beta ;
    A (k:m,k:n) = A (k:m,k:n) - v * (beta * (v' * A (k:m,k:n))) ;
end
R = A ;
```

## MATLAB prototype: left-looking QR

```
function [V,Beta,R] = qr_left (A)
[m n] = size (A) ;
V = zeros (m,n) ;
Beta = zeros (1,n) ;
R = zeros (m,n) ;
for k = 1:n
    x = A (:,k) ;
    for i = 1:k-1
        v = V (i:m,i) ;
        beta = Beta (i) ;
        x (i:m) = x (i:m) - v * (beta * (v' * x (i:m))) ;
    end
    [v,beta,s] = gallery ('house', x (k:m), 2) ;
    V (k:m,k) = v ;
    Beta (k) = beta ;
    R (1:(k-1),k) = x (1:(k-1)) ;
    R (k,k) = s ;
end
```

## Nonzero pattern of $HA$

Theorem (George, Liu, and Ng)

*Consider  $HA = A - v(\beta(v^T A))$ . Then  $(HA)_{i*}$  where  $i \notin \mathcal{V}$  is equal to row  $i$  of  $A$ . For any row  $i \in \mathcal{V}$ , the nonzero pattern of  $(HA)_{i*}$  is*

$$\bigcup_{i \in \mathcal{V}} \mathcal{A}_{i*}. \quad (1)$$

*That is, in  $HA$ , the nonzero pattern of any modified row  $i \in \mathcal{V}$  is replaced with the set union of all rows that are modified by the Householder reflection  $H$ .*

## Nonzero pattern of $R$

Theorem (Golub and Van Loan [?])

*If  $A^T A$  is positive definite, and its Cholesky factorization is  $LL^T = A^T A$ , then  $L = R_{1\dots n, 1\dots n}^T$ .*

## Nonzero pattern of $R$ , continued

Theorem (Coleman et al., George and Heath )

*Assuming the matrix  $A$  has the strong Hall property,  $\mathcal{R}_{*k} = \mathcal{L}_{k*}$ , where  $\mathcal{L}_{k*}$  denotes the nonzero pattern of the  $k$ th row of the symbolic Cholesky factor of  $A^T A$ . If  $A$  does not have the strong Hall property,  $\mathcal{R}_{*k} \subseteq \mathcal{L}_{k*}$ .*



## More concisely ...

### Theorem

$\mathcal{R}_{*k} = \text{Reach}_{\mathcal{T}_k}(\{\min \mathcal{A}_{i*} \mid i \in \mathcal{A}_{*k}\})$  (assuming  $A$  has the strong Hall property).

# Nonzero pattern of the Householder vectors

## Theorem

$$\mathcal{V}_k = \left( \bigcup_{k=\text{parent}(i)} \mathcal{V}_i \setminus \{i\} \right) \cup \{i \mid k = \min \mathcal{A}_{i*}\},$$

where each set in the above expression is disjoint from all other sets, and  $A$  has the strong Hall property. That is,

$$|\mathcal{V}_k| = \left( \sum_{k=\text{parent}(i)} |\mathcal{V}_i| - 1 \right) + |\{i \mid k = \min \mathcal{A}_{i*}\}| \quad (2)$$

If  $A$  does not have the strong Hall property, this is an upper bound on  $\mathcal{V}$ .

## left-looking QR

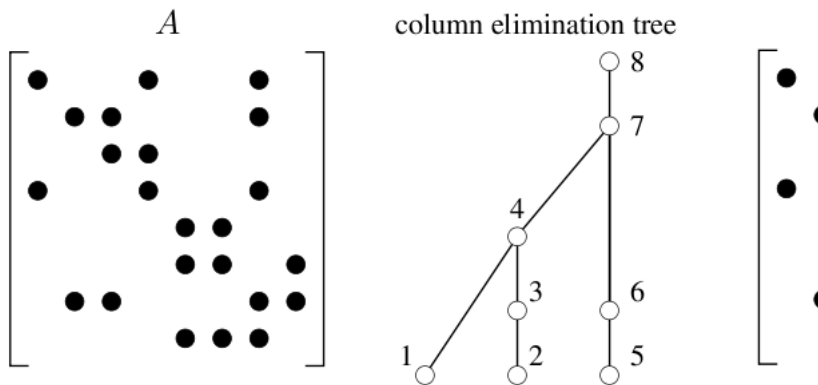


Figure 5.1. *QR factorization*

## left-looking QR

```

function  $[V, \beta_{1\dots n}, R] = \text{sparse\_qr\_left}(A)$ 
     $\mathcal{T} =$  elimination tree of  $A^T A$ 
    compute  $|R|$  using cs_counts of  $A^T A$ 
    compute  $|\mathcal{V}_{1\dots n}|$  using (5.2)
    for  $k = 0$  to  $n - 1$  do
         $\mathcal{R}_{*k} = \text{Reach}_{\mathcal{T}_k}(\{\min \mathcal{A}_{i*} \mid i \in \mathcal{A}_{*k}\})$ 
         $x = A_{*k}$ 
         $\mathcal{V}_k = \mathcal{A}_{*k}$ 
        for each  $i \in \mathcal{R}_{*k}$  do
             $x = x - v_i(\beta_i(v_i^T x))$ 
            if  $\text{parent}(i) = k$  then
                 $\mathcal{V}_k = \mathcal{V}_k \cup \mathcal{V}_i \setminus \{i\}$ 
         $R_{1\dots k-1, k} = x_{1\dots k-1}$ 
         $[v_k, \beta_k, r_{kk}] = \text{house}(x_{k\dots m})$ 

```