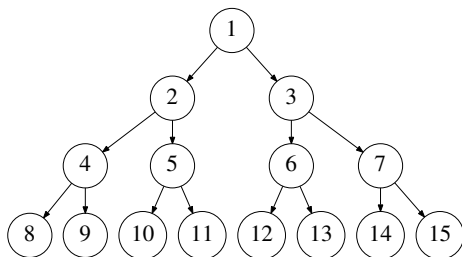


Search and Game Playing

- Uninformed search
- Informed search
- Iterative improvement, Constraint satisfaction
- Game playing

1

Search Problems: Definition



Search = \langle initial state, operators, goal states \rangle

- Initial State: description of the current situation as given in a problem
- Operators: functions from any state to a set of successor (or neighbor) states
- Goal: subset of states, or test rule

3

Overview: Uninformed search

- Search problems: definition
- Example: 8-puzzle
- General search
- Evaluation of search strategies
- Strategies: breadth-first, uniform-cost, depth-first
- More uninformed search: depth-limited, iterative deepening, bidirectional search

2

Variants of Search Problems

Search = \langle state space, initial state, operators, goal states \rangle

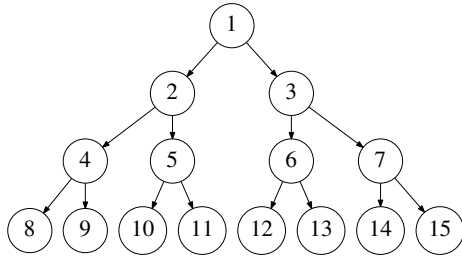
- State space: set of all possible states reachable from the current initial state through repeated application of the operators (i.e. path).

Search = \langle initial state, operators, goal states, path cost \rangle

- Path cost: find **the best** solution, not just **a** solution. Cost can be many different things.

4

Types of Search



- Uninformed: systematic strategies
- Informed: Use domain knowledge to narrow search
- Game playing as search: minimax, state pruning, probabilistic games

5

Operators

Function from state to subset of states

- drive to neighboring city
- place piece on chess board
- add person to meeting schedule
- slide tile in 8-puzzle

Characteristics

- often requires instantiation (fill in variables)
- encode constraints (only certain operations are allowed)
- generally discrete: continuous parameters → infinite branching

7

Search State

State as Data Structure

- examples: variable assignment, properties, order in list, bitmap, graph (vertex and edges)
- captures all possible ways world could be
- typically static, discrete (symbolic), but does not have to be

Choosing a Good Representation

- concise (keep only the relevant features)
- explicit (easy to compute when needed)
- embeds constraints

6

Goals: Subset of states or test rules

Specification:

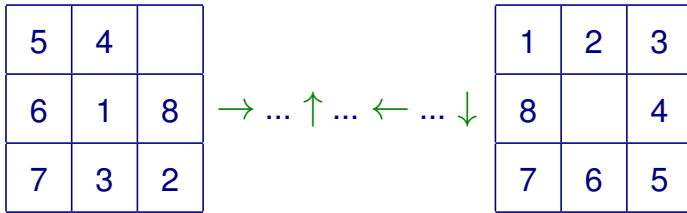
- set of states: enumerate the eligible states
- partial description: e.g. a certain variable has value over x .
- constraints: or set of constraints. Hard to enumerate all states matching the constraints, or very hard to come up with a solution at all (i.e. you can only verify it; P vs. NP).

Other considerations:

- space, time, quality (exact vs. approximate trade-offs)

8

An Example: 8-Puzzle

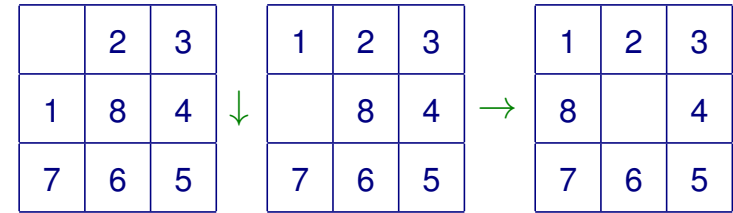


- **State:** location of 8 number tiles and one blank tile
- **Operators:** blank moves left, right, up, or down
- **Goal test:** state matches the configuration on the right (see above)
- **Path cost:** each step cost 1, i.e. path length, or search tree depth

Generalization: 15-puzzle, ..., $(N^2 - 1)$ -puzzle

9

8-Puzzle: Example



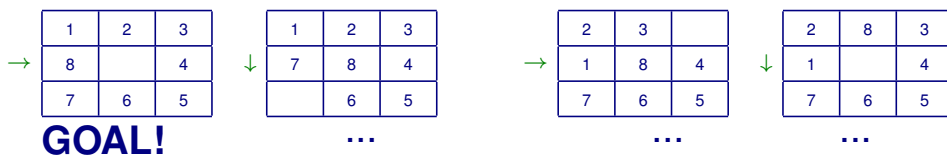
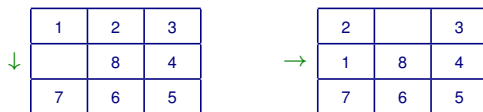
Possible state representations in LISP (0 is the blank):

- (0 2 3 1 8 4 7 6 5)
- ((0 2 3) (1 8 4) (7 6 5))
- ((0 1 7) (2 8 6) (3 4 5))
- or use the `make-array`, `aref` functions.

How easy to: (1) compare, (2) operate on, and (3) store (i.e. size).

10

8-Puzzle: Search Tree



Goal Test

As simple as a single LISP call:

```
* (defvar *goal-state* '(1 2 3 8 0 4 7 6 5))
*GOAL-STATE*

* (equal *goal-state* '(1 2 3 8 0 4 7 6 5))
T
```

11

12

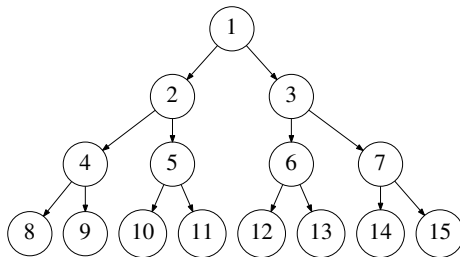
General Search Algorithm

Pseudo-code:

```
function General-Search (problem, Que-Fn)
  node-list := initial-state
  loop begin
    // fail if node-list is empty
    if Empty(node-list) then return FAIL
    // pick a node from node-list
    node := Get-First-Node(node-list)
    // if picked node is a goal node, success!
    if (node == goal) then return as SOLUTION
    // otherwise, expand node and enqueue
    node-list := Que-Fn(node-list, Expand(node))
  loop end
```

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Breadth First Search



- node visit order (goal test): 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
- queuing function: enqueue at end (add expanded node at the end of the list)

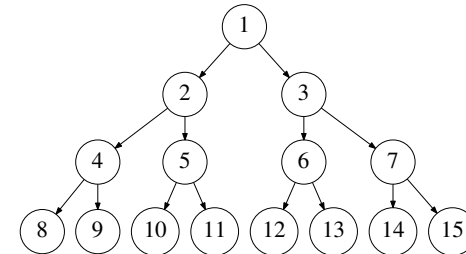
15

Evaluation of Search Strategies

- time-complexity: how many nodes expanded so far?
- space-complexity: how many nodes must be stored in node-list at any given time?
- completeness: if solution exists, guaranteed to be found?
- optimality: guaranteed to find the best solution?

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BFS: Expand Order



Evolution of the queue (**bold**= expanded and added children):

1. **[1]** : initial state
2. **[2][3]** : dequeue 1 and enqueue 2 and 3
3. [3] **[4][5]** : dequeue 2 and enqueue 4 and 5
4. [4] [5] **[6][7]** : **all** depth 3 nodes
- ...
8. [8] [9] [10] [11] [12] [13] **[14][15]** : **all** depth 4 nodes

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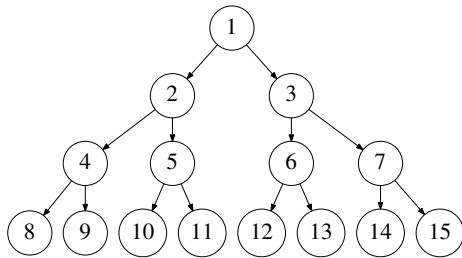
BFS: Evaluation

branching factor b , depth of solution d :

- complete: it will find the solution if it exists
- time: $1 + b + b^2 + \dots + b^d$
- space: $O(b^{d+1})$ where d is the depth of the shallowest solution
- space is more problem than time in most cases (p 75, figure 3.12).
- time is also a major problem nonetheless (same as time)

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Depth First Search



- node visit order (goal test): 1 2 4 8 9 5 10 11 3 6 12 13 7 14 15
- queuing function: enqueue at left (stack push; add expanded node at the beginning of the list)

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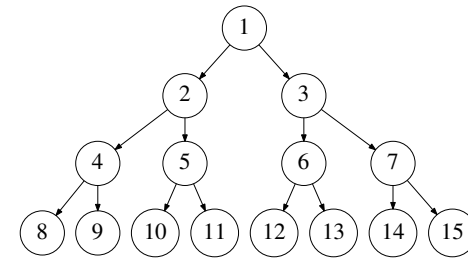
Uniform Cost

BFS with expansion of lowest-cost nodes: path cost is $g(\text{node})$.

- BFS: $g(n) = \text{Depth}(\text{node})$

18

DFS: Expand Order



Evolution of the queue (**bold**=expanded and added children):

1. [1] : initial state
2. **[2]**[3] : pop 1 and push expanded in the front
3. **[4]****[5]** [3] : pop 2 and push expanded in the front
4. **[8]****[9]** [5] [3] : pop 4 and push expanded in the front

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DFS: Evaluation

branching factor b , depth of solutions d , max depth m :

- incomplete: may wander down the wrong path
- time: $O(b^m)$ nodes expanded (worst case)
- space: $O(bm)$ (just along the current path)
- good when there are many shallow goals
- bad for deep or infinite depth state space

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Key Points

- Description of a search problem: initial state, goals, operators, etc.
- Considerations in designing a representation for a state
- Evaluation criteria
- BFS, UCS, DFS: time and space complexity, completeness
- Differences and similarities between BFS and UCS
- When to use one vs. another
- Node visit orders for each strategy
- Tracking the stack or queue at any moment

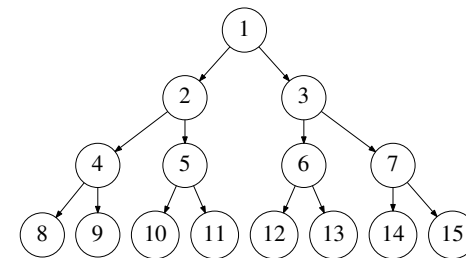
23

Implementation

- Use of stack or queue : explicit storage of expanded nodes
- Recursion : implicit storage in the recursive call stack

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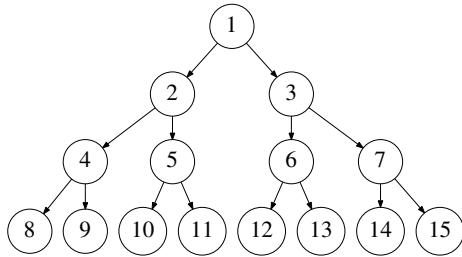
Depth Limited Search (DLS): Limited Depth DFS



- node visit order for each depth limit l :
1 ($l = 1$); 1 2 3 ($l = 2$); 1 2 4 5 3 6 7 ($l = 3$);
- queuing function: enqueue at front (i.e. stack push)
- **push the depth of the node as well:**
($\langle \text{depth} \rangle \langle \text{node} \rangle$)

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DLS: Expand Order



Evolution of the queue (**bold**=expanded and then added):

($\langle \text{depth} \rangle, \langle \text{node} \rangle$); Depth limit = 3

1. [(d1, 1)] : initial state
2. [**(d2,2)**][**(d2,3)**] : pop 1 and push 2 and 3
3. [**(d3,4)**][**(d3,5)**] [(d2, 3)] : pop 2 and push 4 and 5
4. [(d3, 5)] [(d2, 3)] : pop 4, cannot expand it further
5. [(d2, 3)] : pop 5, cannot expand it further
6. [**(d3,6)**][**(d3,7)**] : pop 3, and push 6, 7

...

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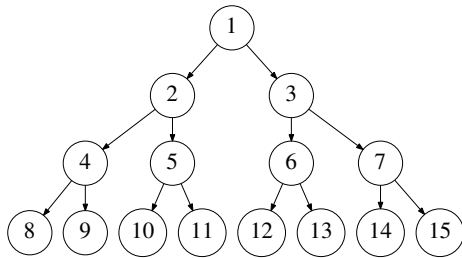
DLS: Evaluation

branching factor b , depth limit l , depth of solution d :

- complete: if $l \geq d$
- time: $O(b^l)$ nodes expanded (worst case)
- space: $O(bl)$ (same as DFS, where $l = m$ (m : max depth of tree in DFS))
- good if solution is within the limited depth.
- non-optimal (same problem as in DFS).

26

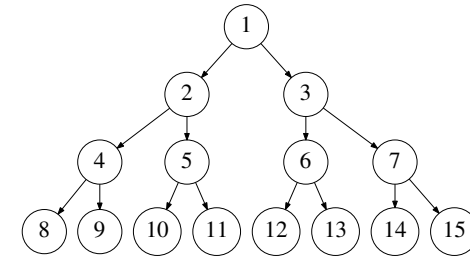
Iterative Deepening Search: DLS by Increasing Limit



- node visit order:
1 ; 1 2 3; 1 2 4 5 3 6 7; 1 2 4 8 9 5 10 11 3 6 12 13 7 14 15; ...
- revisits already explored nodes at successive depth limit
- queuing function: enqueue at front (i.e. stack push)
- **push the depth of the node as well: ($\langle \text{depth} \rangle \langle \text{node} \rangle$)**

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IDS: Expand Order



Basically the same as DLS: Evolution of the queue (**bold**=expanded and then added): ($\langle \text{depth} \rangle, \langle \text{node} \rangle$); e.g. Depth limit = 3

1. [(d1, 1)] : initial state
2. [**(d2,2)**][**(d2,3)**] : pop 1 and push 2 and 3
3. [**(d3,4)**][**(d3,5)**] [(d2, 3)] : pop 2 and push 4 and 5
4. [(d3, 5)] [(d2, 3)] : pop 4, cannot expand it further
5. [(d2, 3)] : pop 5, cannot expand it further
6. [**(d3,6)**][**(d3,7)**] : pop 3, and push 6, 7

...

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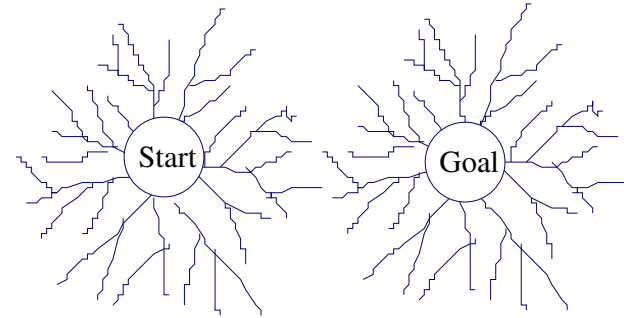
IDS: Evaluation

branching factor b , depth of solution d :

- complete: cf. DLS, which is conditionally complete
- time: $O(b^d)$ nodes expanded (worst case)
- space: $O(bd)$ (cf. DFS and DLS)
- **optimal!:** unlike DFS or DLS
- good when search space is huge and the depth of the solution is not known (*)

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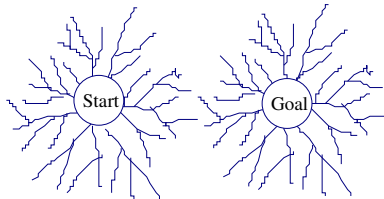
Bidirectional Search (BDS)



- Search from both initial state and goal to reduce search depth.
- $O(b^{d/2})$ of BDS vs. $O(b^{d+1})$ of BFS.

30

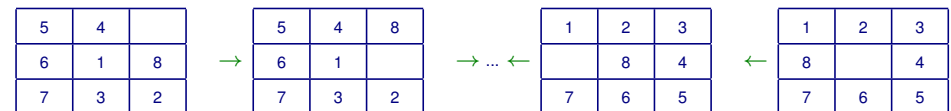
BDS: Considerations



1. how to back trace from the goal?
2. successors and predecessors: are operations reversible?
3. are goals explicit?: need to know the goal to begin with
4. check overlap in two branches
5. BFS? DFS? which strategy to use? Same or different?

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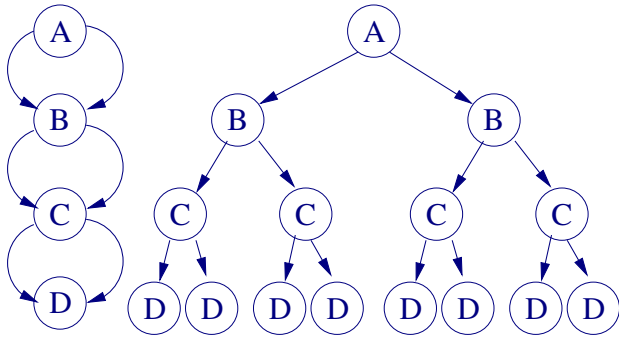
BDS Example: 8-Puzzle



- Is it a good strategy?
- What about Chess? Would it be a good strategy?
- What kind of domains may be suitable for BDS?

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Avoiding Repeated States



Repeated states can be devastating in search problems.

- Common cases: problems with reversible operators → search space becomes infinite
- One approach: find a spanning tree of the graph

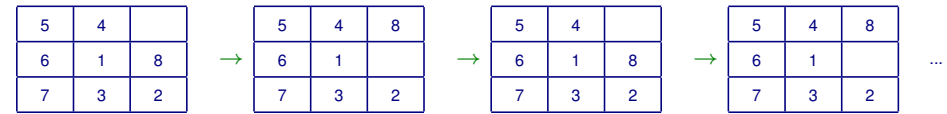
33

Key Points

- DLS, IDS, BDS search order, expansions, and queuing
- DLS, IDS, BDS evaluation
- DLS, IDS, BDS: suitable domains
- Repeated states: why removing them is important

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Avoiding Repeated States: Strategies



- Do not return to the node's parent
- Avoid cycles in the path (this is a huge theoretical problem in its own right)
- Do not generate states that you generated before: use a hash table to make checks efficient

How to avoid storing every state? Would using a short signature (or a checksum) of the full state description help?

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Overview: Informed search, Iterative Improvement

- Best-first search
- Heuristic function
- Greedy best-first search
- A^*
- Designing good heuristics
- IDA^*
- Iterative improvement algorithms
 1. Hill-climbing
 2. Simulated annealing

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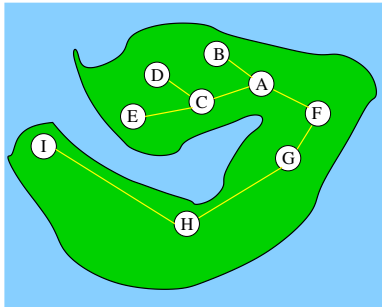
Informed Search

From domain knowledge, obtain an **evaluation function**.

- best-first search: order nodes according to the evaluation function value
- greedy search: minimize estimated cost for reaching the goal – fast, but incomplete and non-optimal.
- A^* : minimize $f(n) = g(n) + h(n)$, where $g(n)$ is the current path cost from start to n , and $h(n)$ is the estimated cost from n to goal.

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Heuristic Function



- $h(n)$ = estimated cost of the cheapest path from the state at node n to a goal state.
- The only requirement is the $h(n) = 0$ at the goal.
- **Heuristics** means “to find” or “to discover”, or more technically, “how to solve problems” (Polya, 1957).

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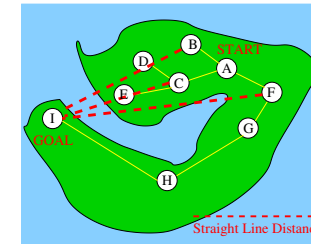
Best First Search

```
function Best-First-Search (problem, Eval-Fn)
    Queuing-Fn ← sorted list by Eval-Fn(node)
    return General-Search(problem, Queuing-Fn)
```

- The queuing function queues the expanded nodes, and sorts it every time by the $Eval-Fn$ value of each node.
- One of the simplest $Eval-Fn$: **estimated cost** to reach the goal.

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Heuristics: Example



- $h_{SLD}(n)$: straight line distance (SLD) is one example.
- Start from **A** and Goal is **I**: **C** is the most promising next step in terms of $h_{SLD}(n)$, i.e. $h(C) < h(B) < h(F)$
- Requires some knowledge:
 1. coordinates of each city
 2. generally, cities toward the goal tend to have smaller **SLD**.

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Greedy Best-First Search

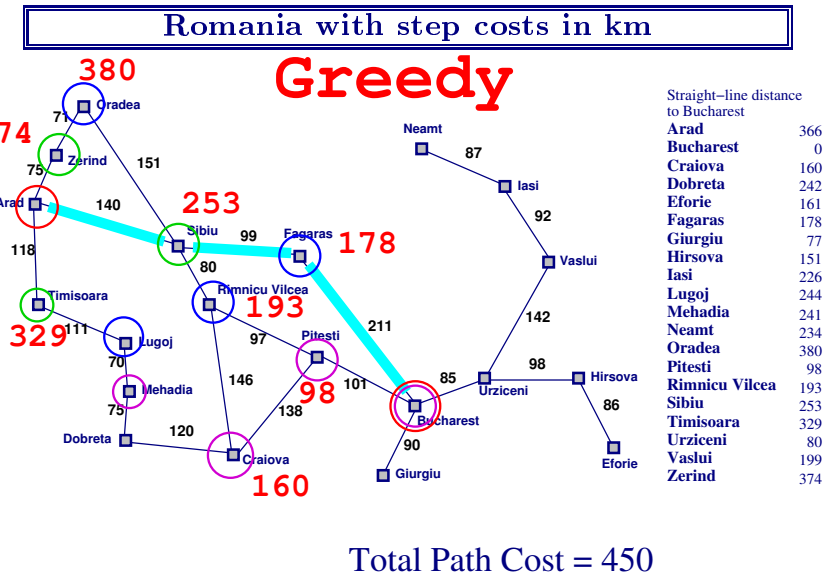
function Greedy-Best-First Search (*problem*)

$h(n)$ = estimated cost from n to goal

return Best-First-Search(*problem*, h)

- Best-first with heuristic function $h(n)$

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AIMA Slides ©Stuart Russell and Peter Norvig, 1998

Chapter 4, Sections 1-2, 4 5

Greedy Best-First Search: Evaluation

Branching factor b and max depth m :

- Fast, just like Depth-First-Search: single path toward the goal.
- Time: $O(b^m)$
- Space: same as time – all nodes are stored in sorted list(!), **unlike DFS**
- Incomplete, just like DFS
- Non-optimal, just like DFS

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A^* : Uniform Cost + Heuristic Search

Avoid expanding paths that are already found to be expensive:

- $f(n) = g(n) + h(n)$
- $f(n)$: estimated cost to goal through node n
- **provably complete and optimal!**
- **restrictions:** $h(n)$ should be an **admissible heuristic**
- admissible heuristic: one that **never overestimate** the actual cost of the best solution through n

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A* Search

function A^* -Search (*problem*)

$g(n)$ =current cost up till n

$h(n)$ =estimated cost from n to goal

return Best-First-Search(*problem*, $g + h$)

- Condition: $h(n)$ must be an **admissible heuristic function!**
- A^* is **optimal!**

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Behavior of A* Search

- usually, the f value never decreases along a given path:
monotonicity
- in case it is nonmonotonic, i.e. $f(Child) < f(Parent)$, make this adjustment:
 $f(Child) = \max(f(Parent), g(Child) + h(Child))$.
- this is called **pathmax**

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Optimality of A*

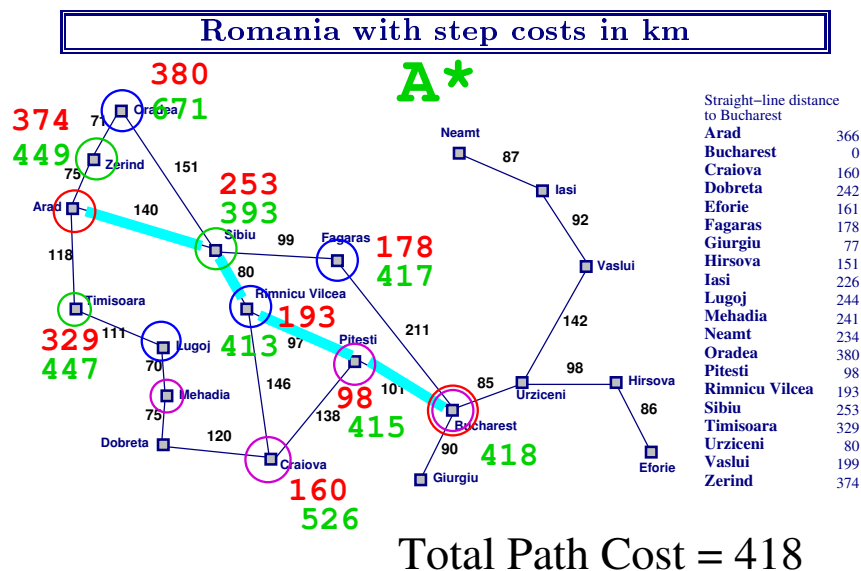
G_2 : suboptimal goal in the node-list.

n : unexpanded node on a shortest path to goal G_1

- $f(G_2) = g(G_2)$ since $h(G_2) = 0$
- $> g(G_1)$ since G_2 is suboptimal
- $\geq f(n)$ since h is admissible

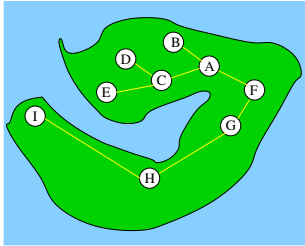
Since $f(G_2) > f(n)$, A^* will never select G_2 for expansion.

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Optimality of A^* : Example



1. **Expansion of parent allowed:** search fails at nodes **B**, **D**, and **E**.
2. **Expansion of parent disallowed:** paths through nodes **B**, **D**, and **E** will have an inflated path cost $g(n)$, thus will become nonoptimal.

$A \rightarrow C \rightarrow E \rightarrow C \rightarrow A \rightarrow F \rightarrow \dots$
 inflated path cost

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Complexity of A^*

A^* is complete and optimal, but space complexity can become exponential if the heuristic is not good enough.

- condition for **subexponential** growth:

$$|h(n) - h^*(n)| \leq O(\log h^*(n)),$$

where $h^*(n)$ is the **true** cost from n to the goal.

- that is, error in the estimated cost to reach the goal should be less than even linear, i.e. $< O(h^*(n))$.

Unfortunately, with most heuristics, error is at least proportional with the true cost, i.e. $\geq O(h^*(n)) > O(\log h^*(n))$.

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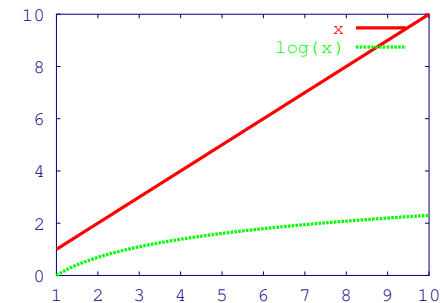
Lemma to Optimality of A^*

Lemma: A^* expands nodes in order of increasing $f(n)$ value.

- Gradually adds **f-contours** of nodes (cf. BFS adds layers).
- The goal state may have a f value: let's call it f^*
- This means that all nodes with $f < f^*$ will be expanded!

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Linear vs. Logarithmic Growth Error



- Error in heuristic: $|h(n) - h^*(n)|$.
- For most heuristics, the error is at least linear.
- For A^* to have subexponential growth, the error in the heuristic should be on the order of $O(\log h^*(n))$.

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Problem with A^*

Space complexity is usually **exponential!**

- we need a memory bounded version
- one solution is: Iterative Deepening A^* , or IDA^*

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Heuristic Functions: Example

Eight puzzle

5	4		1	2	3
6	1	8	8		4
7	3	2	7	6	5

- $h_1(n)$ = number of misplaced tiles
- $h_2(n)$ = total **Manhattan** distance (city block distance)

$$h_1(n) = 7 \text{ (not counting the blank tile)}$$

$$h_2(n) = 2+3+3+2+4+2+0+2 = 18$$

* Both are admissible heuristic functions.

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A^* : Evaluation

- Complete : unless there are infinitely many nodes with $f(n) \leq f(G)$
- Time complexity: exponential in (relative error in $h \times$ length of solution)
- Space complexity: same as time (keep all nodes in memory)
- Optimal

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Dominance

If $h_2(n) \geq h_1(n)$ for all n and both are admissible, then we say that $h_2(n)$ **dominates** $h_1(n)$, and is better for search.

Typical search costs for depth $d = 14$:

- Iterative Deepening : 3,473,941 nodes expanded
- $A^*(h_1)$: 539 nodes
- $A^*(h_2)$: 113 nodes

Observe that in A^* , every node with $f < f^*$ is expanded. Since $f = g + h$, nodes with $h(n) < f^* - g(n)$ will be expanded, so larger h will result in less nodes being expanded.

- f^* is the f value for the optimal solution path.

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Designing Admissible Heuristics

Relax the problem to obtain an admissible heuristics.

For example, in 8-puzzle:

- allow tiles to move anywhere $\rightarrow h_1(n)$
- allow tiles to move to any adjacent location $\rightarrow h_2(n)$

For traveling:

- allow traveler to travel by air, not just by road: **SLD**

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Iterative Deepening A^* : IDA^*

A^* is complete and optimal, but the performance is limited by the available space.

- Basic idea: only search within a certain f bound, and gradually increase the f bound until a solution is found.
- More on IDA^* next time.

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Other Heuristic Design

- Use composite heuristics: $h(n) = \max(h_1(n), \dots, h_m(n))$
- Use statistical information: random sample h and true cost to reach goal. Find out how often h and true cost is related.

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IDA^*

```
function  $IDA^*$ (problem)
    root  $\leftarrow$  Make-Node(Initial-State(problem))
    f-limit  $\leftarrow$  f-Cost(root)
    loop do
        solution, f-limit  $\leftarrow$  DFS-Contour(root, f-limit)
        if solution != NULL then return solution
        if f-limit ==  $\infty$  then return failure
    end loop
```

Basically, iterative deepening depth-first-search with depth defined as the f -cost ($f = g + n$):

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DFS-Contour(*root*, *f-limit*)

Find solution from node **root**, within the *f*-cost limit of **f-limit**.

DFS-Contour returns **solution sequence** and new *f*-cost limit.

- if $f\text{-cost}(\text{root}) > \text{f-limit}$, return fail.
- if **root** is a goal node, return solution and new *f*-cost limit.
- **recursive call** on all successors and return solution and **minimum *f*-limit** returned by the calls
- return **null solution** and new *f*-limit by default

Similar to the recursive implementation of DFS.

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IDA* : Time Complexity

Depends on the heuristics:

- small number of possible heuristic function values \rightarrow small number of *f*-contours to explore \rightarrow becomes similar to A*
- complex problems: each *f*-contour only contain one new node
if A* expands N nodes,
IDA* expands
$$1 + 2 + \dots + N = \frac{N(N+1)}{2} = O(N^2)$$
- a possible solution is to have a **fixed** increment ϵ for the *f*-limit
 \rightarrow solution will be suboptimal for at most ϵ (ϵ -admissible)

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IDA* : Evaluation

- complete and optimal (with same restrictions as in A*)
- space: proportional to longest path that it explores (because it is depth first!)
- time: dependent on the number of different values $h(n)$ can assume.

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Other Methods: Beam Search

Best-first search with a fixed limited branching factor

- expand the first n nodes with the best Eval-Fn value, where n is a small number.
- n is called the **width of the beam**
- good for domains with continuous time functions (like speech recognition)
- good for domains with huge branching factor (like above)

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Iterative Improvement Algorithms

Start with a complete configuration (all variable values assigned, and **optimal**), and **gradually improve** it.

- Hill-climbing (maximize cost function)
- Gradient descent (minimize cost function)
- Simulated Annealing (probabilistic)

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Hill-Climbing Strategies

Problems of local maxima, plateau, and ridges:

- try **random-restart**: move to a random location in the landscape and restart search from there
- keep n best nodes (beam search) *
- parallel search
- simulated annealing *

Hardness of problem depends on the **shape of the landscape**.

*: coming up next

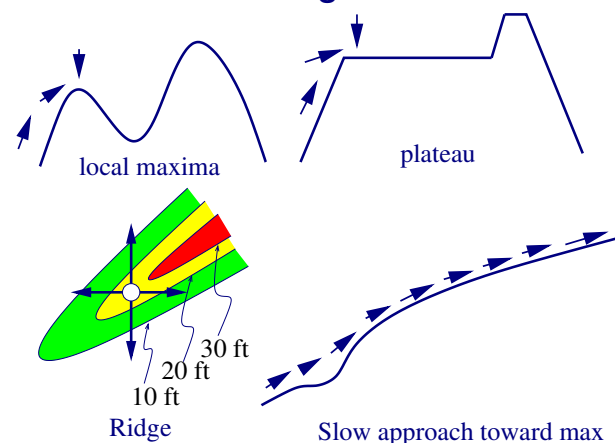
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Hill-Climbing

- no queue, keep only the best node
- greedy, no back-tracking
- good for domains where **all nodes are solutions**:
 - goal is to **improve** quality of the solution
 - optimization problems
- note that it is different from greedy search, which keeps a node list

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Hill-Climbing: Problems



- Possible solution: **simulated annealing** – gradually decrease randomness of move to attain globally optimal solution (more on this next week).

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Simulated Annealing: Overview

Annealing:

- heating metal to a high-temperature (making it a liquid) and then allowing to cool slowly (into a solid); this relieves internal stresses and results in a more stable, lower-energy state in the solid.
- at high temperature, atoms move actively (large distances with greater randomness), but as temperature is lowered, they become more static.

Simulated annealing is similar:

- basically, hill-climbing with randomness that allows going **down** as well as the standard **up**
- randomness (as temperature) is reduced over time

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Simulated Annealing (SA)

Goal: minimize the energy E , as in statistical thermodynamics.

For successors of the current node,

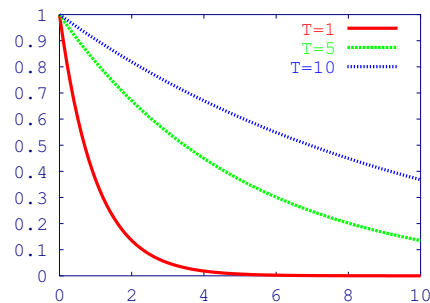
- if $\Delta E \leq 0$, the move is accepted
- if $\Delta E > 0$, the move is accepted with probability $P(\Delta E) = e^{-\frac{\Delta E}{kT}}$, where k is the Boltzmann constant and T is temperature.
- randomness is in the comparison: $P(\Delta E) < \text{rand}(0, 1)$

$$\Delta E = E_{\text{new}} - E_{\text{old}}.$$

The heuristic $h(n)$ or $f(n)$ represents E .

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Temperature and $P(\Delta E) < \text{rand}(0, 1)$



Downward moves of any size are allowed at high temperature, but at low temperature, only small downward moves are allowed.

- Higher temperature $T \rightarrow$ higher probability of **downward** hill-climbing
- Lower $\Delta E \rightarrow$ higher probability of **downward** hill-climbing

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T Reduction Schedule

High to low temperature reduction schedule is important:

- reduction too fast: suboptimal solution
- reduction too slow: wasted time
- question: does the form of the reduction schedule curve matter?
linear, quadratic, exponential, etc.?

The proper values are usually found experimentally.

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Simulated Annealing Applications

- VLSI wire routing and placement
- Various scheduling optimization tasks
- Traffic control
- Neural network training
- etc.

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Constraints

- Unary, binary, and higher order constraints: how many variables should simultaneously meet the constraint
- Absolute constraints vs. preference constraints
- Variables are defined in a certain **domain**, which determines the possible set of values, either discrete or continuous.

This is part of a much more complex problem called **constrained optimization problems** in operations research consisting of cost function (either minimize or maximize) and several constraints.

Problems can be linear, nonlinear, convex, nonconvex, etc.

Straight-forward solutions exist for a limited subclass of these (for example, for linear programming problems can be solved by the simplex method).

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Constraint Satisfaction Search

Constraint Satisfaction Problem (CSP):

- **state**: values of a set of **variables**
- **goal**: test if a set of constraints are met
- **operators**: set values of variables
- general search can be used, but specialized solvers for CSP work better

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CSP: continued

- CSPs include NP-complete problems such as 3-SAT, thus finding the solutions can require exponential time.
- However, constraints can help narrow down the possible options, therefore reducing the branching factor. This is because in CSP, the goal can be decomposed into several constraints, rather than being a whole solution.
- Strategies: backtracking (back up when constraint is violated), forward checking (do not expand further if look-ahead returns a constraint violation). Forward checking is often faster and simple to implement.

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Heuristics for Constraint Satisfaction Problems

General strategies for **variable selection**:

- Most-constrained-variable heuristic (var with fewest possible values)
- Most-constraining-variable heuristic (var involved in the largest number of constraints)

and for **value assignment**:

- Least-constraining-value heuristic (value that rules out the smallest number of values for vars)

Reducing branching factor vs. leaving freedom for future choices.

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Game Playing

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Key Points

- best-first-search: definition
- heuristic function $h(n)$: what it is
- greedy search: relation to $h(n)$ and evaluation. How it is different from DFS (time complexity, space complexity)
- A^* : definition, evaluation, conditions of optimality
- complexity of A^* : relation to error in heuristics
- designing good heuristics: several rule-of-thumbs
- IDA^* : evaluation, time and space complexity (worst case)
- beam search concept
- hill-climbing concept and strategies
- simulated annealing: core algorithm, effect of T and ΔE , source of randomness.

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Game Playing

- attractive AI problem because it is **abstract**
- one of the oldest domains in AI
- in most cases, the world state is fully accessible
- computer representation of the situation can be clear and exact
- challenging: uncertainty introduced by the opponent and the complexity of the problem (full search is impossible)
- hard: in chess, branching factor is about 35, and 50 moves by each player = 35^{100} nodes to search
 - compare to 10^{40} possible legal board states
- *game playing is more like real life than mechanical search*

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Games vs. Search Problems

“Unpredictable” opponent → solution is a contingency plan

Time limits → unlikely to find goal, must approximate

Plan of attack:

- algorithm for perfect play (Von Neumann, 1944)
- finite horizon, approximate evaluation (Zuse, 1945; Shannon, 1950; Samuel, 1952–57)
- pruning to reduce costs (McCarthy, 1956)

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Types of Games

	deterministic	chance
perfect info	chess, checkers, go, othello	backgammon, monopoly
imperfect info	?	bridge, poker, scrabble

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Two-Person Perfect Information Game

initial state: initial position and who goes first

operators: legal moves

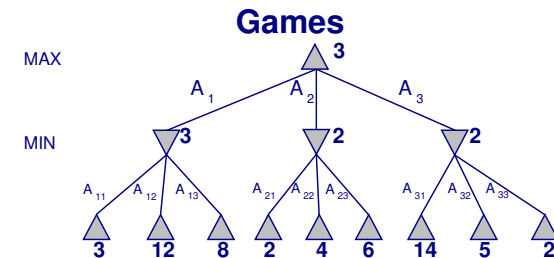
terminal test: game over?

utility function: outcome (win:+1, lose:-1, draw:0, etc.)

- two players (**MIN** and **MAX**) taking turns to maximize their chances of winning (each turn generates one **ply**)
- one player's victory is another's defeat
- need a **strategy** to win no matter what the opponent does

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Minimax: Strategy for Two-Person Perfect Info Games



- generate the whole tree, and apply util function to the leaves
- go back upward assigning utility value to each node
- at MIN node, assign **min(successors' utility)**
- at MAX node, assign **max(successors' utility)**
- **assumption:** the opponent acts optimally

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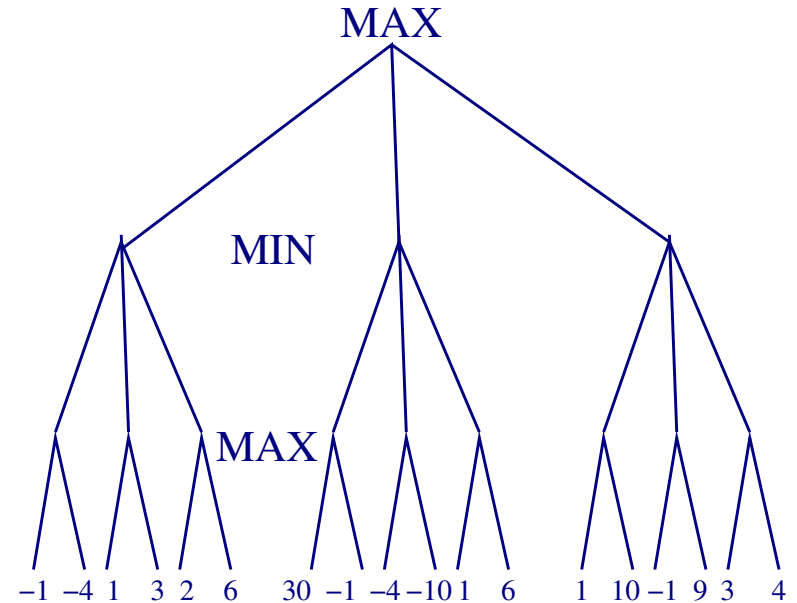
Minimax Decision

```
function Minimax-Decision (game) returns operator
  return operator that leads to a child state with the
  max(Minimax-Value(child state,game))
```

```
function Minimax-Value(state,game) returns utility value
  if Goal(state), return Utility(state)
  else if Max's move then
    → return max of successors' Minimax-Value
  else
    → return min of successors' Minimax-Value
```

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Minimax Exercise



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Minimax: Evaluation

Branching factor b , max depth m :

- **complete**: if the game tree is finite
- **optimal**: if opponent is optimal
- **time**: b^m
- **space**: bm – depth-first (only when utility function values of all nodes are known!)

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Resource Limits

- Time limit: as in Chess → can only evaluate a fixed number of paths
- Approaches:
 - **evaluation function** : how desirable is a given state?
 - **cutoff test** : depth limit
 - **pruning**

Depth limit can result in the **horizon effect**: interesting or devastating events can be just over the horizon!

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Evaluation Functions

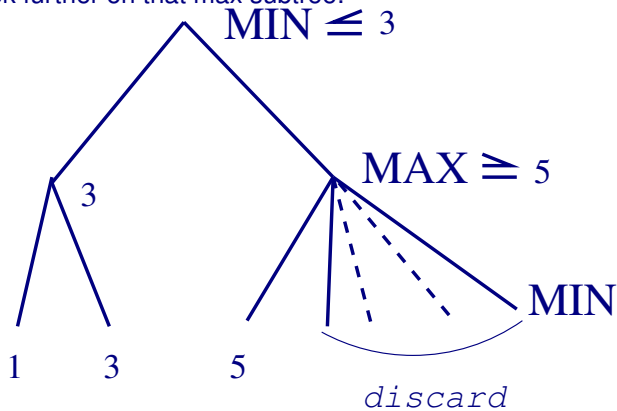
For chess, usually a **linear** weighted sum of feature values:

- $Eval(s) = \sum_i w_i f_i(s)$
- $f_i(s) = (\text{number of white piece X}) - (\text{number of black piece X})$
- other features: degree of control over the center area
- exact values do not matter: the **order** of Minimax-Value of the successors matter.

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β Cuts

When the current min value is less than the successor's max value, don't look further on that max subtree:

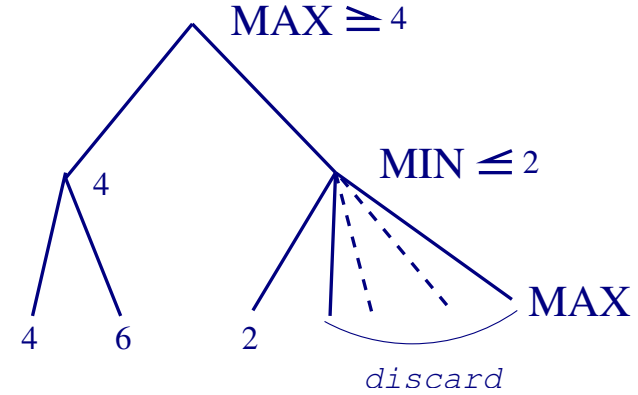


Right subtree can be **at least 5**, so MIN will always choose the left path regardless of what appears next.

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α Cuts

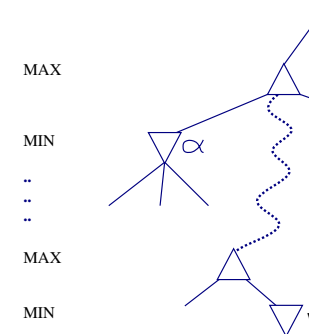
When the current max value is greater than the successor's min value, don't look further on that min subtree:



Right subtree can be **at most 2**, so MAX will always choose the left path regardless of what appears next.

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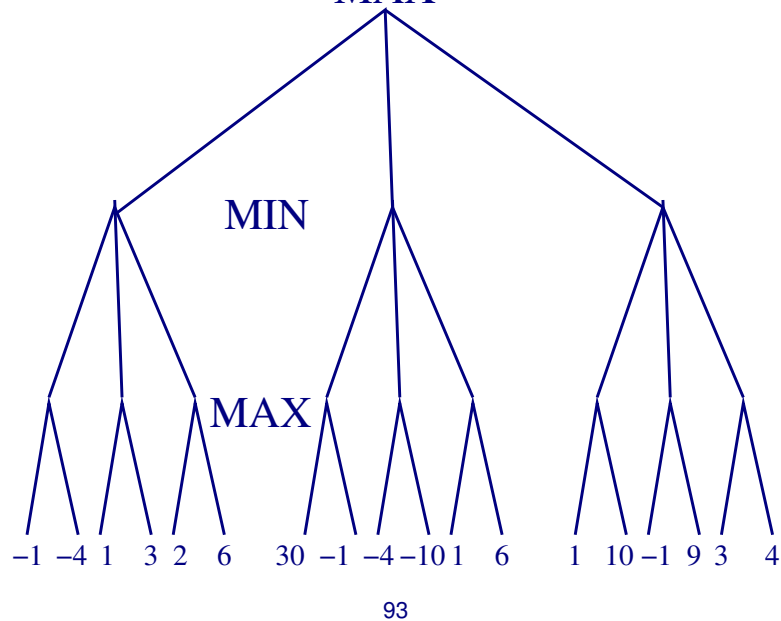
$\alpha - \beta$ Pruning



- memory of best MAX value α and best MIN value β
- do not go further on any one that does worse than the remembered α and β

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$\alpha - \beta$ Exercise MAX



Key Points

- Game playing: what are the types of games?
- Minimax: definition, and how to get minmax values
- Minimax: evaluation
- α - β pruning: why it saves time

$\alpha - \beta$ Pruning Properties

Cut off nodes that are known to be suboptimal.

Properties:

- pruning **does not** affect final result
- good move ordering improves effectiveness of pruning
- with **perfect ordering**, time complexity = $b^{m/2}$
 - **doubles** depth of search
 - can easily reach 8-ply in chess
- $b^{m/2} = (\sqrt{b})^m$, thus $b = 35$ in chess reduces to $b = \sqrt{35} \approx 6$!!!

Up-coming Topics

- formal $\alpha - \beta$ pruning algorithm
- $\alpha - \beta$ pruning properties
- games with an element of chance
- state-of-the-art game playing with AI
- more complex games

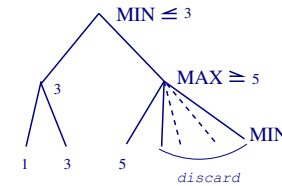
$\alpha - \beta$ Pruning: Initialization

Along the path from the beginning to the current **state**:

- α : best MAX value
 - initialize to $-\infty$
- β : best MIN value
 - initialize to ∞

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$\alpha - \beta$ Pruning Algorithm: Max-Value

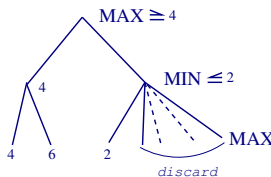


```

function Max-Value (state, game,  $\alpha$ ,  $\beta$ ) return utility value
 $\alpha$ : best MAX on path to state ;  $\beta$ : best MIN on path to state
if Cutoff(state) then return Utility(state)
 $v \leftarrow -\infty$ 
for each s in Successor(state) do
  ·  $v \leftarrow \text{Max}(\alpha, \text{Min-Value}(s, \text{game}, \alpha, \beta))$ 
  · if  $v \geq \beta$  then return v /* CUT!! */
  ·  $\alpha \leftarrow \text{Max}(\alpha, v)$ 
end
return v
    
```

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$\alpha - \beta$ Pruning Algorithm: Min-Value



```

function Min-Value (state, game,  $\alpha$ ,  $\beta$ ) return utility value
 $\alpha$ : best MAX on path to state ;  $\beta$ : best MIN on path to state
if Cutoff(state) then return Utility (state)
 $v \leftarrow \infty$ 
for each s in Successor(state) do
  ·  $v \leftarrow \text{Min}(\beta, \text{Max-Value}(s, \text{game}, \alpha, \beta))$ 
  · if  $v \leq \alpha$  then return v /* CUT!! */
  ·  $\beta \leftarrow \text{Min}(\beta, v)$ 
end
return v
    
```

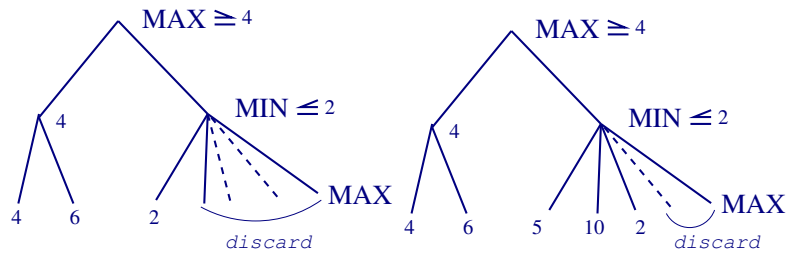
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$\alpha - \beta$ Pruning Tips

- At a MAX node:
 - Only α is updated with the MAX of successors.
 - Cut is done by checking if returned $v \geq \beta$.
 - If all fails, MAX(v of successors) is returned.
- At a MIN node:
 - Only β is updated with the MIN of successors.
 - Cut is done by checking if returned $v \leq \alpha$.
 - If all fails, MIN(v of successors) is returned.

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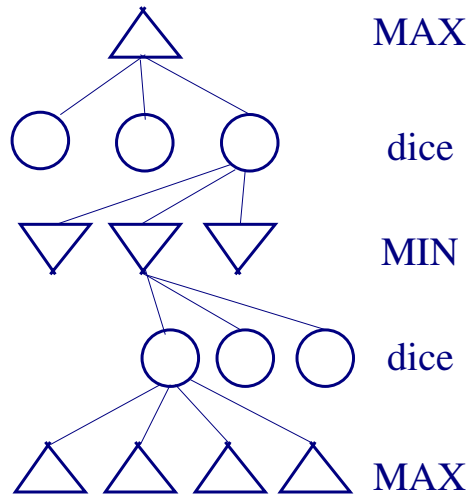
Ordering is Important for Good Pruning



- For MIN, sorting successor's utility in an **increasing** order is better (shown above; left).
- For MAX, sorting in **decreasing** order is better.

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Game Tree With Chance Element



- chance element forms a new ply (e.g. dice, shown above)

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Games With an Element of Chance

Rolling the dice, shuffling the deck of card and drawing, etc.

- **chance nodes** need to be included in the minimax tree
- try to make a move that maximizes the **expected value** → **expectimax**
- expected value of random variable X :

$$E(X) = \sum_x xP(x)$$

- expectimax

$$\text{expectimax}(C) = \sum_i P(d_i) \max_{s \in S(C, d_i)} (\text{utility}(s))$$

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Design Considerations for Probabilistic Games

- the **value** of evaluation function, not just the **scale** matters now! (think of what expected value is)
- time complexity: $b^m n^m$, where n is the number of distinct dice rolls
- pruning can be done if we are careful

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Historical Achievements in AI Game Playing

- Checkers: Samuel's Checker Program running on 10Kbyte (1952)
- Backgammon: Tesauro's Neural Network → top three (1992)
- Chess: IBM's Deep Blue defeated Garry Kasparov (1997)
- Othello (Reversi): smaller search space → superhuman performance

Genetic algorithms can perform very well on select domains.

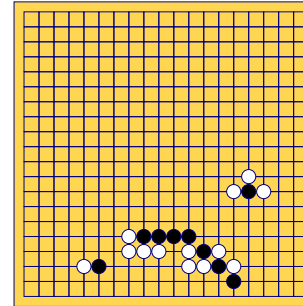
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Google DeepMind's AlphaGo

- AlphaGo by Google DeepMind beat the Korean Go grand master Lee Sedol in 2016 (4 to 1).
- AlphaGo combined deep reinforcement learning with tree search.

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Hard Games



The game of *Go*, popular in East Asia:

- $19 \times 19 = 361$ grid: branching factor is huge!
- search methods inevitably fail: need more structured rules
- the bet was high: \$1,400,000 prize for the first computer program to beat a select, 12-year old player. The late Mr. Ing Chang Ki (photo above) put up the money from his personal funds.

Photo from <http://www.samsloan.com/ing.htm>.

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Deep Learning based Game Playing

- Atari 2600 games: DeepMind (2015)
- Chess, Shogi (Japanese chess), Go: Alpha Zero (2017)
- Atari games + Multiple board games, without knowledge of the rules: MuZero (2019)
- Starcraft II : AlphaStar (2019)

Google DeepMind really seems to like games! :-)

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Key Points

- formal $\alpha - \beta$ pruning algorithm: know how to apply pruning
- $\alpha - \beta$ pruning properties: evaluation
- games with an element of chance: what are the added elements?
how does the minmax tree get augmented?

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Emacs Tips

M-x : [Alt]-[x] or [ESC] then [x], **C-x**: [CTRL]-[x]

- M-x shell (run shell within emacs)
- C-p (\uparrow), C-n (\downarrow), C-b (\leftarrow), C-f (\rightarrow)
- C-x C-f (load file)
- M-x lisp-mode (environment for editing lisp code)
- C-s (search forward) C-r (reverse search)
- C-g (abort current command in scratch)
- C-k (cut line) C-y (yank, or paste)
- C-space (begin block) C-x C-x (end block) C-w (cut) C-y (yank, or paste)
- C-x u or M-x undo (undo) ; C-x C-s (save) ; C-x C-c (exit)

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