Basics of Probability Theory

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The **probability space** or **sample space** Ω is the set of all possible outcomes of an experiment. For example, the sample space of the coin tossing experiment is $\Omega = \{\text{head}, \text{tail}\}.$

Certain subsets of the sample space are called **events**, and the probability of these events is determined by a **probability measure**.

If we roll a dice, then one of its six face values is the outcome of the experiment, so the sample space is $\Omega = \{1, 2, 3, 4, 5, 6\}$.

An event is a subset of the sample space Ω . The event $\{1,2\}$ occurs when the dice shows a face value less than three.

The probability measures describes the odds that a certain event occurs, for instance $\Pr[\{1,2\}] = 1/3$ means that the event $\{1,2\}$ will occur with probability 1/3.

A probability measure is not necessarily defined on all subsets of the sample space Ω , but just on all subsets of Ω that are considered events. Nevertheless, we want to have a uniform way to reason about the probability of events. This is accomplished by requiring that the collection of events form a σ -algebra.

A σ -algebra \mathcal{F} is a collection of subsets of the sample space Ω such that the following requirements are satisfied:

S1 The empty set is contained in \mathcal{F} .

S2 If a set *E* is contained in \mathcal{F} , then its complement E^c is contained in \mathcal{F} .

S3 The countable union of sets in \mathcal{F} is contained in \mathcal{F} .

The empty set \emptyset is often called the **impossible event**.

The sample space Ω is the complement of the empty set, hence is contained in \mathcal{F} . The event Ω is called the **certain event**.

If *E* is an event, then $E^c = \Omega \setminus E = \Omega - E$ is called the **complementary event**.

Let \mathcal{F} be a σ -algebra.

Exercise If A and B are events in \mathcal{F} , then $A \cap B$ in \mathcal{F} .

Exercise

The countable intersection of events in \mathcal{F} is contained in \mathcal{F} .

Exercise

If A and B are events in \mathcal{F} , then $A - B = A \setminus B$ is contained in \mathcal{F} .

Example

Remark

Let \mathcal{A} be a subset of $P(\Omega)$. Then the intersection of all σ -algebras containing \mathcal{A} is a σ -algebra, called the smallest σ -algebra generated by \mathcal{A} . We denote the smallest σ -algebra generated by \mathcal{A} by $\sigma(\mathcal{A})$.

Example

Let
$$\Omega = \{1, 2, 3, 4, 5, 6\}$$
 and $\mathcal{A} = \{\{1, 2\}, \{2, 3\}\}.$

$$\sigma(\mathcal{A}) = \{\emptyset, \{1, 2, 3, 4, 5, 6\}, \{1, 2\}, \{3, 4, 5, 6\}, \{2, 3\}, \{1, 4, 5, 6\}, \{2, 3\}, \{1, 4, 5, 6\}, \{1\}, \{2, 3, 4, 5, 6\}, \{1\}, \{2, 3, 4, 5, 6\}, \{1, 2, 3\}, \{4, 5, 6\}, \{1, 3\}, \{2, 4, 5, 6\}\}$$

Exercise

Let $\Omega = \{1, 2, 3, 4, 5, 6\}$ and $\mathcal{A} = \{\{2\}, \{1, 2, 3\}, \{4, 5\}\}$. Determine $\sigma(\mathcal{A})$.

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Determine $\sigma(\mathcal{A})$.

Solution

We have

$$\begin{aligned} \mathcal{A} &= \{ \varnothing, \Omega, \{2\}, \{1, 3, 4, 5, 6\}, \\ & \{1, 2, 3\}, \{4, 5, 6\}, \{4, 5\}, \{1, 2, 3, 6\}, \\ & \{1, 3\}, \{2, 4, 5, 6\}, \{6\}, \{1, 2, 3, 4, 5\}, \\ & \{2, 6\}, \{1, 3, 4, 5\}, \{2, 4, 5\}, \{1, 3, 6\} \} \end{aligned}$$

Let \mathcal{F} be a σ -algebra over the sample space Ω . A **probability measure** on \mathcal{F} is a function $\Pr: \mathcal{F} \to [0, 1]$ satisfying

- **P1** The certain event satisfies $Pr[\Omega] = 1$.
- **P2** If the events E_1, E_2, \ldots in \mathcal{F} are mutually disjoint, then

$$\Pr\left[\bigcup_{k=1}^{\infty} E_k\right] = \sum_{k=1}^{\infty} \Pr[E_k].$$

Example

Example

Probability Function Let Ω be a sample space and let $a \in \Omega$. Suppose that $\mathcal{F} = P(\Omega)$ is the σ -algebra. Then Pr: $\Omega \to [0, 1]$ given by

$$\mathsf{Pr}[\mathcal{A}] = egin{cases} 1 & ext{ if } \mathbf{a} \in \mathcal{A}, \ 0 & ext{ otherwise}. \end{cases}$$

is a probability measure.

We know that **P1** holds, since $Pr[\Omega] = 1$. **P2** holds as well. Indeed, if E_1, E_2, \ldots are mutually disjoint events in $P(\Omega)$, then at most one of the events contains *a*.

$$\sum_{k=1}^{\infty} \Pr[E_k] = \begin{cases} 1 & \text{if some set } E_k \text{ contains } a, \\ 0 & \text{if none of the sets } E_k \text{ contains } a. \end{cases} = \Pr[\bigcup_{k=1}^{\infty} E_k]$$

These axioms have a number of familiar consequences. For example, it follows that the complementary event E^c has probability

$$\Pr[E^c] = 1 - \Pr[E].$$

In particular, the impossible event has probability zero, $Pr[\emptyset] = 0$.

Another consequence is a simple form of the **inclusion-exclusion principle**:

$$\Pr[E \cup F] = \Pr[E] + \Pr[F] - \Pr[E \cap F],$$

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which is convenient when calculating probabilities. Indeed,

$$Pr[E \cup F] = Pr[E \setminus (E \cap F)] + Pr[E \cap F] + Pr[F \setminus (E \cap F)]$$

= Pr[E] + Pr[F \ (E \circ F)] + (Pr[E \circ F] - Pr[E \circ F])
= Pr[E] + Pr[F] - Pr[E \circ F].

Exercise

Let E and F be events such that $E \subseteq F$. Show that

 $\Pr[E] \leq \Pr[F].$

Exercise

Let E_1, \ldots, E_n be events that are not necessarily disjoint. Show that $\Pr[E_1 \cup \cdots \cup E_n] \leq \Pr[E_1] + \cdots + \Pr[E_n].$

Conditional Probabilities

Let *E* and *F* be events over a sample space Ω such that $\Pr[F] > 0$. The **conditional probability** $\Pr[E|F]$ of the event *E* given *F* is defined by

$$\Pr[E|F] = \frac{\Pr[E \cap F]}{\Pr[F]}.$$

The value $\Pr[E|F]$ is interpreted as the probability that the event *E* occurs, assuming that the event *F* occurs.

By definition, $Pr[E \cap F] = Pr[E|F] Pr[F]$, and this simple multiplication formula often turns out to be useful.

Law of Total Probability (Simplest Version)

Law of Total Probability

Let Ω be a sample space and A and E events. We have

$$Pr[A] = Pr[A \cap E] + Pr[A \cap E^{c}]$$

= Pr[A | E] Pr[E] + Pr[A | E^{c}] Pr[E^{c}].

The events *E* and *E^c* are disjoint and satisfy $\Omega = E \cup E^c$. Therefore, we have

$$\Pr[A] = \Pr[A \cap E] + \Pr[A \cap E^c].$$

The second equality follows directly from the definition of conditional probability.

Bayes' Theorem (Simplest Version)

Bayes' Theorem

$$\Pr[A \mid B] = \frac{\Pr[B \mid A] \Pr[A]}{\Pr[B]}$$

We have

$$\Pr[A \mid B] \Pr[B] = \Pr[A \cap B] = \Pr[B \cap A] = \Pr[B \mid A] \Pr[A].$$

Dividing by Pr[B] yields the claim.

Bayes' Theorem (Second Version)

Bayes' Theorem (Version 2)

$$Pr[A \mid B] = \frac{Pr[B \mid A] Pr[A]}{Pr[B|A] Pr[A] + Pr[B|A^{c}] Pr[A^{c}]}.$$

By the first version of Bayes' theorem, we have

$$\Pr[A \mid B] = \frac{\Pr[B \mid A] \Pr[A]}{\Pr[B]}$$

Now apply the law of total probability with $\Omega = A \cup A^c$ to the probability $\Pr[B]$ denominator.

Polynomial Identities

Suppose that we use a library that is supposedly implementing a polynomial factorization. We would like to check whether the polynomials such as

$$p(x) = (x+1)(x-2)(x+3)(x-4)(x+5)(x-6)$$

$$q(x) = x^6 - 7x^3 + 25$$

are the same.

We can multiply the terms both polynomials and simplify. This uses $\Omega(d^2)$ multiplications for polynomials of degree d.

If the polynomials p(x) and q(x) are the same, then we must have

 $p(x)-q(x)\equiv 0.$

If the polynomials p(x) and q(x) are not the same, then an integer $r \in \mathbf{Z}$ such that

$$p(r)-q(r)\neq 0$$

would be a **witness** to the difference of p(x) and q(x).

We can check whether $r \in \mathbf{Z}$ is a witness in O(d) multiplications.

We get the following randomized algorithm for checking whether p(x) and q(x) are the same.

Input: Two polynomials p(x) and q(x) of degree d. for i = 1 to t do r = random(1..100d); return 'different' if $p(r) - q(r) \neq 0$ end return 'same' If $p(x) \equiv q(x)$, then every $r \in \mathbf{Z}$ is a non-witness.

If $p(x) \neq q(x)$, then an integer r in the range $1 \leq r \leq 100d$ is a witness if and only if it is not a root of p(x) - q(x). The polynomial p(x) - q(x) has at most d roots.

The probability that the algorithm will return 'same' when the polynomials are different is at most

$$\Pr['same'|p(x) \neq q(x)] \leqslant \left(\frac{d}{100d}\right)^t = \frac{1}{100^t}.$$

Independent Events

Definition Two events *E* and *F* are called **independent** if and only if $Pr[E \cap F] = Pr[E]Pr[F].$

Two events that are not independent are called **dependent**.

Example

Suppose that we flip a fair coin twice. Then the sample space is $\{HH, HT, TH, TT\}$. The probability of each elementary event is given by 1/4. For instance, $Pr[\{HH\}] = 1/4$.

The event *E* that the **first coin is heads** is given by $\{HH, HT\}$. We have Pr[E] = 1/2. The event *F* that **the second coin is tails** is given by $\{HT, TT\}$. We have Pr[F] = 1/2.

Then $E \cap F$ models the event that **the first coin is heads and the second coin is tails**. The events *E* and *F* are independent, since

$$\Pr[E \cap F] = \frac{1}{4} = \Pr[E]\Pr[F].$$

If E and F are independent, then

$$\Pr[E \mid F] = \frac{\Pr[E \cap F]}{\Pr[F]} = \frac{\Pr[E]\Pr[F]}{\Pr[F]} = \Pr[E].$$

In this case, whether or not F happened has no bearing on the probability of E.

Suppose that $E_1, E_2, ..., E_n$ are events. The events are called **mutually independent** if and only if for all subsets *S* of $\{1, 2, ..., n\}$, we have

$$\Pr\left[\bigcap_{i\in S} E_i\right] = \prod_{i\in S} \Pr[E_i].$$

Please note that it is not sufficient to show this condition for $S = \{1, 2, ..., n\}$, but we really need to show this for all subsets.

Example

We toss a fair coin three times. Consider the events:

- E_1 = the first two values are the same,
- E_2 = the first and last value are the same,

 E_3 = the last two values are the same.

The probabilities are $\Pr[E_1] = \Pr[E_2] = \Pr[E_3] = 1/2$. We have $\Pr[E_1 \cap E_2] = \Pr[E_2 \cap E_3] = \Pr[E_1 \cap E_3] = \Pr[\{HHH, TTT\}] = \frac{1}{4}$. Thus, all three pairs of events are independent. But $\Pr[E_1 \cap E_2 \cap E_3] = \frac{1}{4} \neq \Pr[E_1]\Pr[E_2]\Pr[E_3] = \frac{1}{8}$,

so they are not mutually independent.

Example

A school offers as electives A = athletics, B = band, and C = Mandarin Chinese.

$$\begin{aligned} & \Pr[A \cap B \cap C] = 0.04 \quad \Pr[\overline{A} \cap B \cap C] = 0.2 \\ & \Pr[A \cap B \cap \overline{C}] = 0.06 \quad \Pr[\overline{A} \cap B \cap \overline{C}] = 0.1 \\ & \Pr[A \cap \overline{B} \cap C] = 0.1 \quad \Pr[\overline{A} \cap \overline{B} \cap C] = 0.16 \\ & \Pr[A \cap \overline{B} \cap \overline{C}] = 0 \quad \Pr[\overline{A} \cap \overline{B} \cap \overline{C}] = 0.34 \end{aligned}$$

Then $\Pr[A \cap B \cap C] = 0.04 = \Pr[A] \Pr[B] \Pr[C] = 0.2 \cdot 0.4 \cdot 0.5$. But no two of the three events are pair-wise independent:

$$\Pr[A \cap B] = 0.1 \neq \Pr[A] \Pr[B] = 0.2 \cdot 0.4 = 0.08$$

Verifying Matrix Multiplication

The Problem Let A, B, and C be $n \times n$ matrices over $\mathbf{F}_2 = \mathbf{Z}/2\mathbf{Z}$. Is AB = C?

If we use traditional matrix multiplication, then forming the product of A and B requires $\Theta(n^3)$ scalar operations. Using the fastest known matrix multiplications takes about $\Theta(n^{2.37})$ scalar operations. Can we do better using a randomized algorithm?

A witness for $AB \neq C$ would be a vector v such that

 $ABv \neq Cv$.

We can check whether a vector is a witness in $O(n^2)$ time.

Theorem

If $AB \neq C$, and we choose a vector v uniformly at random from $\{0,1\}^n$, then v is a witness for $AB \neq C$ with probability $\geq 1/2$. In other words,

$$\Pr_{\boldsymbol{\mathcal{V}}\in\mathbf{F}_2^n}[AB\boldsymbol{\mathcal{V}}=C\boldsymbol{\mathcal{V}}\mid AB\neq C]\leqslant \frac{1}{2}.$$

Simple Observation

Lemma

Choosing $v = (v_1, v_2, ..., v_n) \in \mathbf{F}_2^n$ uniformly at random is equivalent to choosing each v_k independently and uniformly at random from \mathbf{F}_2 .

Proof.

If we choose each component v_k independently and uniformly at random from \mathbf{F}_2 , then each vector v in \mathbf{F}_2^n is created with probability $1/2^n$.

Conversely, if $v \in \mathbf{F}_2^n$ is chosen uniformly at random, then the components are independent and $v_k = 1$ with probability 1/2.

Let $D = AB - C \neq 0$. Then ABv = Cv if and only if Dv = 0.

Since $D \neq 0$, the matrix D must have a nonzero entry. Without loss of generality, suppose that $d_{11} \neq 0$.

If Dv = 0, then we must have

$$\sum_{k=1}^n d_{1k}v_k=0.$$

Since $d_{11} \neq 0$, this is equivalent to

$$v_1 = -rac{\sum_{k=2}^n d_{1k} v_k}{d_{11}}.$$

Idea (Principle of Deferred Decisions)

Rather than arguing with the vector $v \in \mathbf{F}_2^n$, we can choose each component of v uniformly at random from \mathbf{F}_2 in order form v_n down to v_1 .

Suppose that the components $v_n, v_{n-1}, \ldots, v_2$ have been chosen. This determines the right-hand side of

$$v_1 = -rac{\sum_{k=2}^n d_{1k} v_k}{d_{11}}.$$

Now there is just one choice of v_1 that will make the equality true, so the probability that this equation is satisfied is at most 1/2. In other words, the probability

$$\Pr[ABv = Cv \mid AB \neq C] \leq 1/2$$