Problem Set 1

Due dates: Electronic submission of .pdf files of this homework is due on 1/24/2019 before 11:00am on ecampus, a signed paper copy of the pdf file is due on 1/24/2019 at the beginning of class.

Name: (put your name here)

Resources. (All people, books, articles, web pages, etc. that have been consulted when producing your answers to this homework)

On my honor, as an Aggie, I have neither given nor received any unauthorized aid on any portion of the academic work included in this assignment. Furthermore, I have disclosed all resources (people, books, web sites, etc.) that have been used to prepare this homework.

Signature: _

Problem 1 (20 points for Homework 1A). Submit your solution to Homework 1A on ecampus. See ecampus for instructions. This part of the homework is due on Tuesday, January 22, 8:00pm.

The problems of Homework 1B follow.

Problem 2 (20 points). Give a self-contained proof of the fact that

$$\log_2(n!) \in \Theta(n \log n).$$

[For part of your argument, you can use results that were given in the lecture, but you should write up the proof in your own words.]

Problem 3 (10 points). Amelia attempted to solve n algorithmic problems. She wrote down one problem per page in her journal and marked the page with e when she was unable to solve the problem and with e when she was able to solve it. So the pages of her journal look like this:



Use the decision tree method to show that any algorithm to find a page with an $\stackrel{\textcircled{\mbox{\scriptsize e}}}{=}$ smiley on has to look at all n pages in the worst case.

Problem 4 (10 points). Amelia attempted to solve n algorithmic problems. She wrote down one problem per page in her journal and marked the page with e when she was unable to solve the problem and with e when she was able to solve it. So the pages of her journal look like this:



Use an adversary method to show that any method to find a page with an \bigotimes smiley on it might have to look at all n pages.

Problem 5 (20 points). Amelia attempted to solve n algorithmic problems, where n is an odd number. She wrote down one problem per page in her journal and marked the page with e when she was unable to solve the problem and with e when she was able to solve it. Suppose that we want to find the pattern e, where she was unable to solve a problem, but was able to solve the subsequent problem.

Find an algorithm that always looks at fewer than n pages but is able to correctly find the pattern when it exists. [Hint: First look at all even pages.]

Problem 6 (20 points). Suppose that we are given a sorted array A[1..n] of n numbers. Our goal is to determine whether or not the array A contains duplicate elements. We will limit ourselves to algorithms that use only the spaceship operator <=> for comparisons, where

```
a <=> b :=
if a < b then return -1
if a = b then return 0
if a > b then return 1
if a and b are not comparable then return nil
```

No other methods will be used to compare or inspect elements of the array.

- (a) Give an efficient (optimal) comparison-based algorithm that decides whether A[1..n] contains duplicates using the spaceship operator for comparisons.
- (b) Use an adversarial argument to show that no algorithm can solve the problem with fewer calls to the comparison operator <=> than the algorithm that you gave in (a).

Checklist:

- $\hfill\square$ Did you add your name?
- □ Did you disclose all resources that you have used? (This includes all people, books, websites, etc. that you have consulted)
- \Box Did you sign that you followed the Aggie honor code?
- \Box Did you solve all problems?
- \Box Did you write the solution in your own words?
- \Box Did you submit the pdf file of your homework?
- \Box Did you submit a signed hardcopy of the pdf file in class?