

Depth-First Search

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Depth-First Search

```
Input:  $G = (V, E)$   
for each node  $u$  in  $V$  do  
    mark  $u$  as unvisited  
od;  
for each unvisited node  $u$  do  
    recursiveDFS( $u$ );  
od;
```

```
recursiveDFS( $u$ ):  
    mark  $u$  as visited;  
    for each unvisited neighbor  $v$  of  $u$  do  
        recursiveDFS( $v$ )  
    od
```

Purpose of this loop:
Create DFS forest. Graph can be
directed or undirected.

DFS Forest

By keeping track of parents, we want to construct a forest resulting from the DFS traversal.

Depth-First Search

```
Input:  $G = (V, E)$   
for each node  $u$  in  $V$  do  
     $parent[u] = nil;$   
    mark  $u$  as unvisited  
od;  
  
for each unvisited node  $u$  do  
     $parent[u] = u;$   
    recursiveDFS( $u$ );  
od;
```

```
recursiveDFS( $u$ ):  
    mark  $u$  as visited;  
    for each unvisited neighbor  $v$  of  $u$  do  
         $parent[v] = u;$  recursiveDFS( $v$ )  
    od
```

Refining DFS

Let us keep track of some interesting information for each node. We will timestamp the steps and record the

- **discovery time**, when the recursive call starts
- **finish time**, when its recursive call ends

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```
Input:  $G = (V, E)$ 
for each node  $u$  in  $V$  do
     $parent[u] = nil$ ;
    mark  $u$  as unvisited
od;

 $time = 0$ ;
for each unvisited node  $u$  do
     $parent[u] = u$ ;
    recursiveDFS( $u$ );
od;
```

```
recursiveDFS( $u$ ):
    mark  $u$  as visited;
     $disc[u] = ++time$ ;
    for each unvisited neighbor  $v$  of  $u$  do
         $parent[v] = u$ ; recursiveDFS( $v$ )
    od;
     $fin[u] = ++time$ ;
```

Running Time of DFS

The first for-loop for initialization takes $O(V)$ time.

The second for-loop in non-recursive wrapper considers each node, so $O(V)$ iterations.

One recursive call is made for each node. In a recursive call for the node u , all its neighbors are checked; so the total time in all recursive calls is $O(E)$.

Total time is $O(V+E)$.

Nested Intervals

Define $[\text{disc}[v], \text{fin}[v]]$ to be the **interval** for node v .

Claim: For any two nodes, either one interval precedes the other or one is enclosed in the other

Indeed, the recursive calls are nested.

Corollary: v is a descendant of u in the DFS forest iff the interval of v is inside the interval of u .

Classifying Edges

Consider edge (u,v) in a directed graph $G = (V,E)$ with respect to its DFS forest

- **tree edge**: v is a child of u
- **back edge**: v is an ancestor of u
- **forward edge**: v is a descendant of u but not a child
- **cross edge**: none of the above

Example of Classifying Edges

